

# **Social Network Analysis**

## **Basic Concepts, Methods & Theory**

University of Cologne  
Johannes Putzke

# Agenda

- Introduction
- Basic Concepts
- Mathematical Notation
- Network Statistics

# Textbooks

- Hanneman & Riddle (2005) *Introduction to Social Network Methods*, available at <http://faculty.ucr.edu/~hanneman/nettext/>
- Wasserman & Faust (1994): *Social Network Analysis – Methods and Applications*, Cambridge: Cambridge University Press.

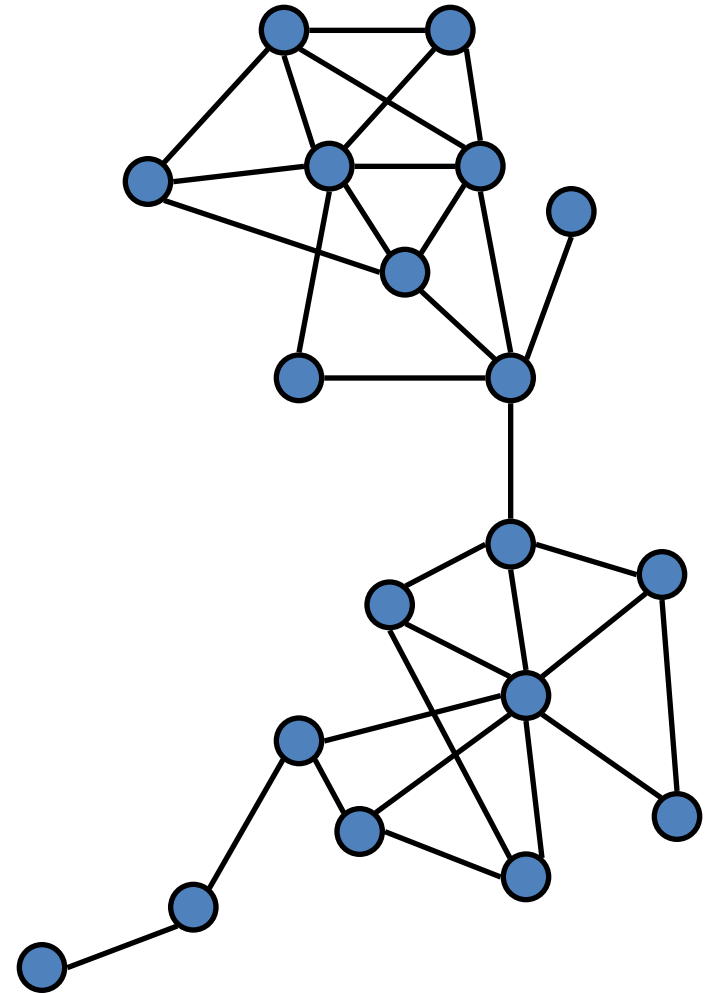
# Introduction

# Basic Concepts

What is a network?

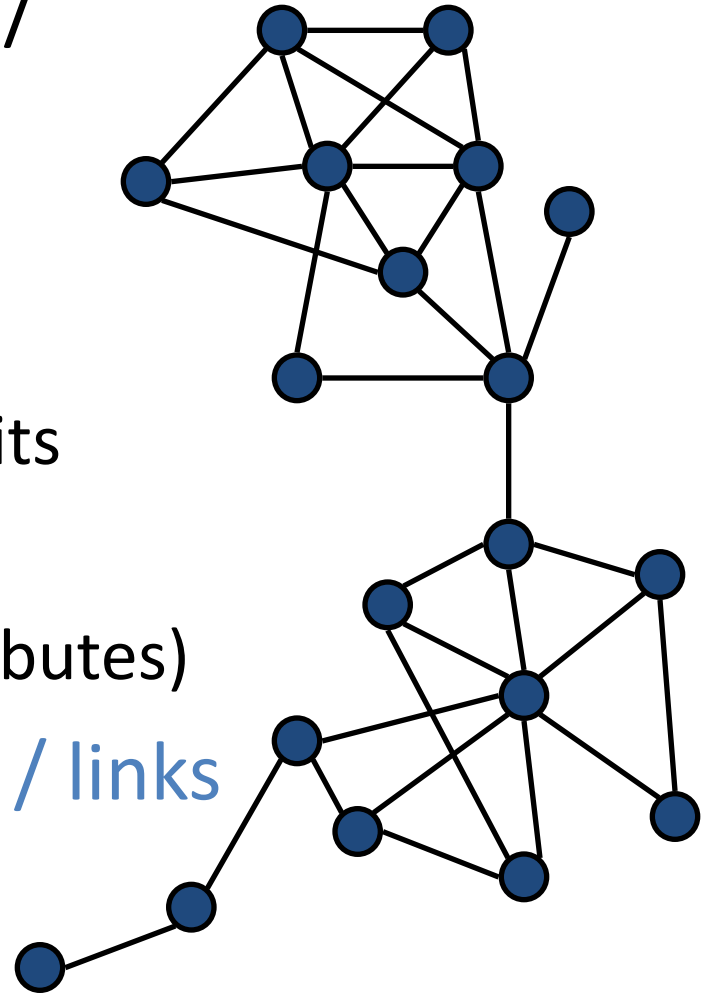
# What is a Network?

- Actors / nodes / vertices / points
- Ties / edges / arcs / lines / links



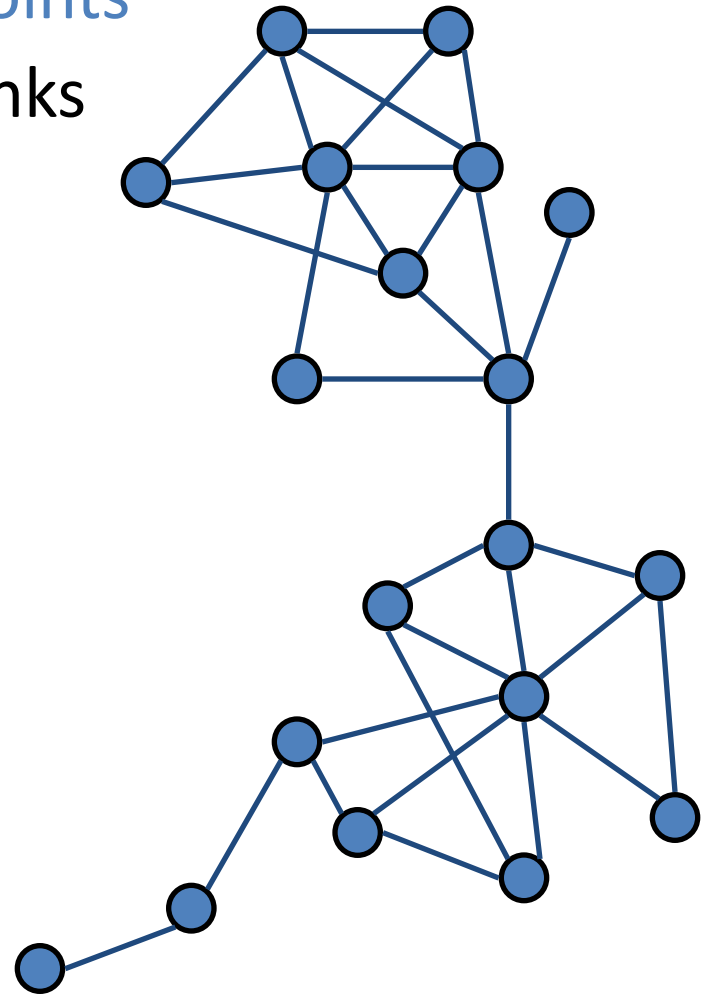
# What is a Network?

- Actors / nodes / vertices / points
  - Computers / Telephones
  - Persons / Employees
  - Companies / Business Units
  - Articles / Books
  - Can have properties (attributes)
- Ties / edges / arcs / lines / links



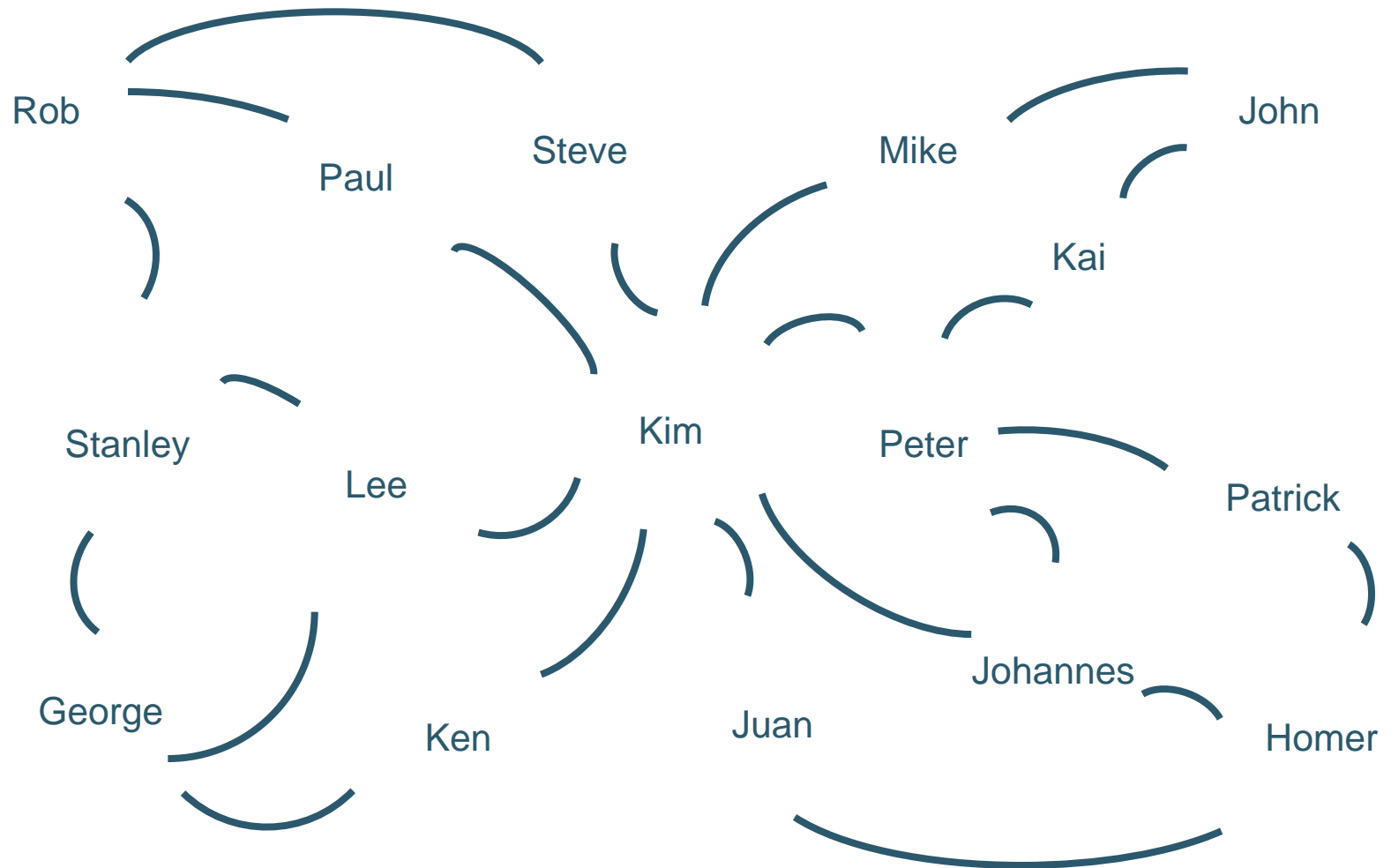
# What is a Network?

- Actors / nodes / vertices / points
- Ties / edges / arcs / lines / links
  - connect pair of actors
  - types of social relations
    - friendship
    - acquaintance
    - kinship
    - advice
    - hindrance
    - sex
  - allow different kind of flows
    - messages
    - money
    - diseases

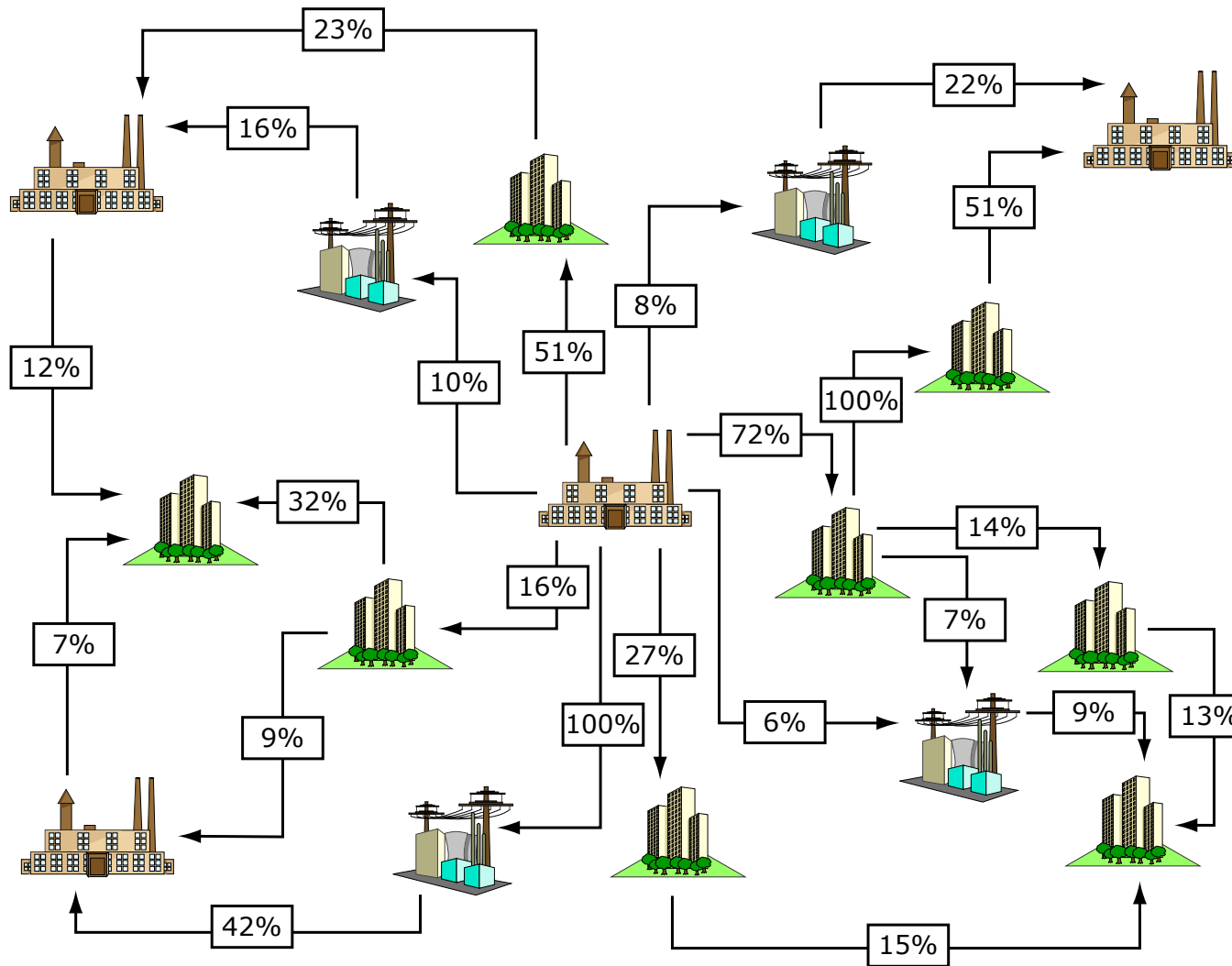




# What is a Social Network? - Relations among People



# What is a Network? - Relations among Institutions



- as institutions
  - owned by, have partnership / joint venture
  - purchases from, sells to
  - competes with, supports
  
- through stakeholders
  - board interlocks
  - Previously worked for

Image by MIT OpenCourseWare.

# Why study social networks?

# Example 2) Homophily Theory

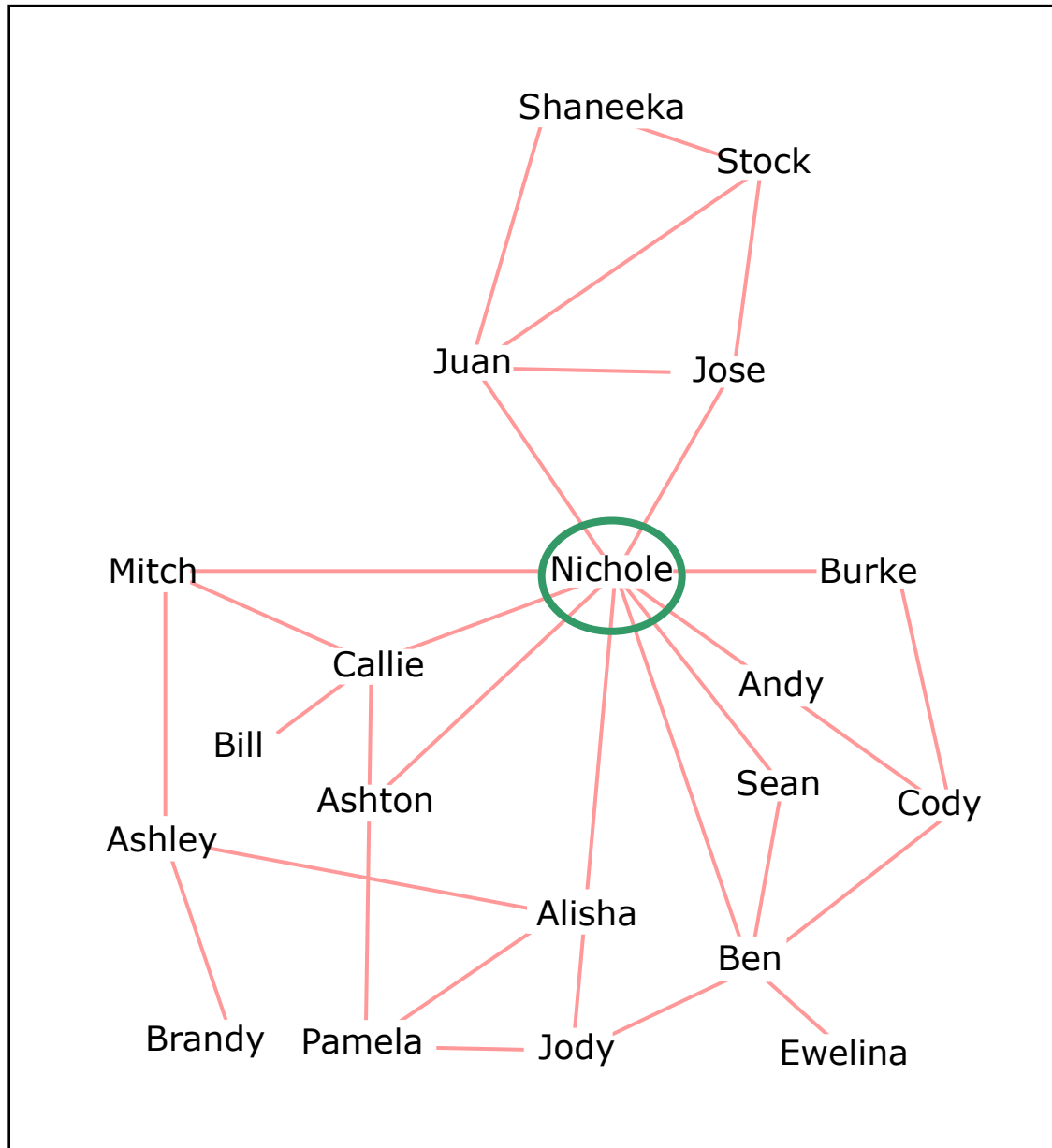
	Male	Female
Male	123	68
Female	95	164

- Birds of a feather flock together
- See McPherson, Smith-Lovin & Cook (2001)

	0-13	14-29	30-44	45-65	>65
0-13	212	63	117	72	91
14-29	83	372	75	67	84
30-44	105	98	321	214	117
45-65	62	72	232	412	148
>65	90	77	124	153	366

- age / gender → network

# Managerial Relevance – Social Network...



# ...vs. Organigram

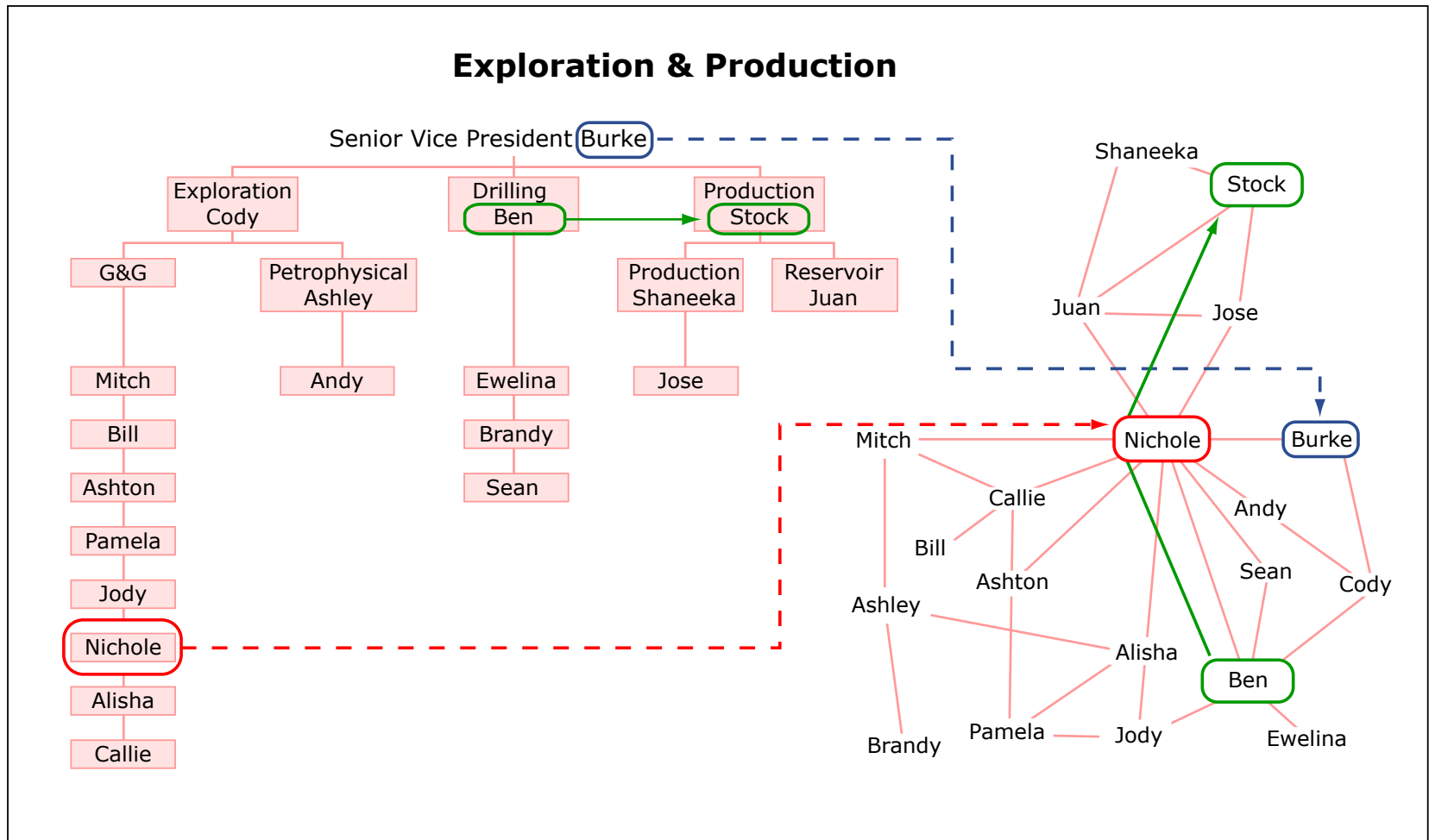
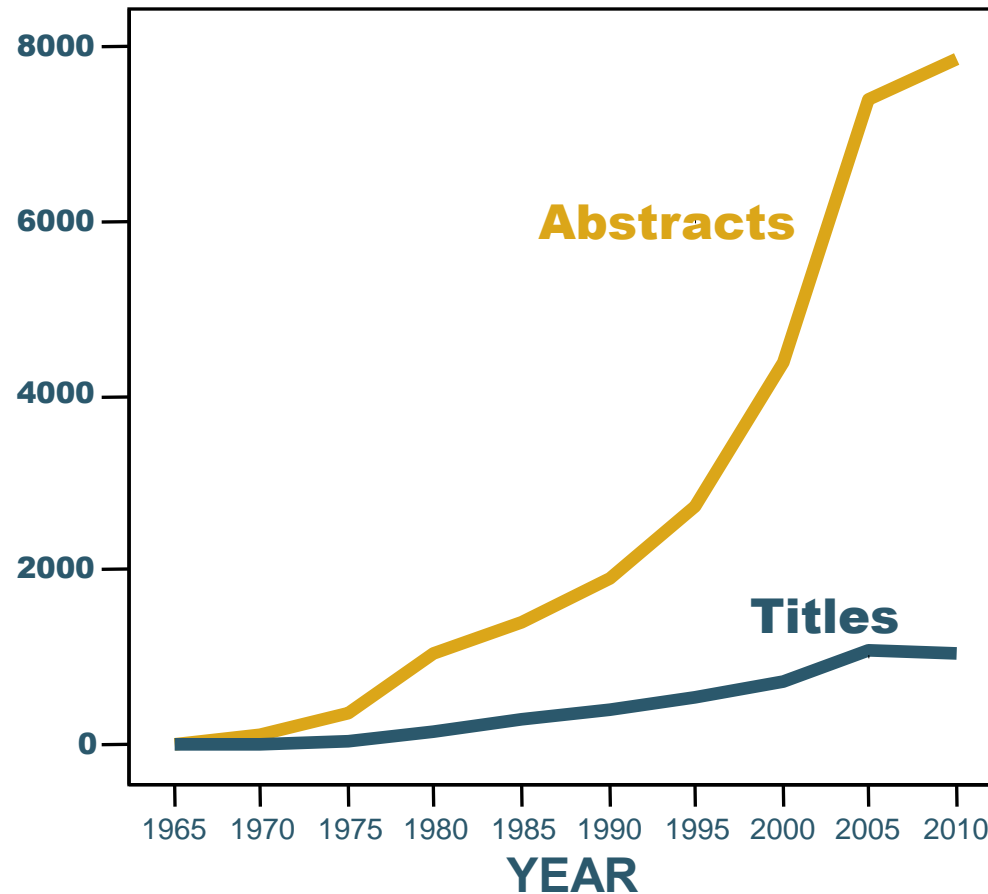


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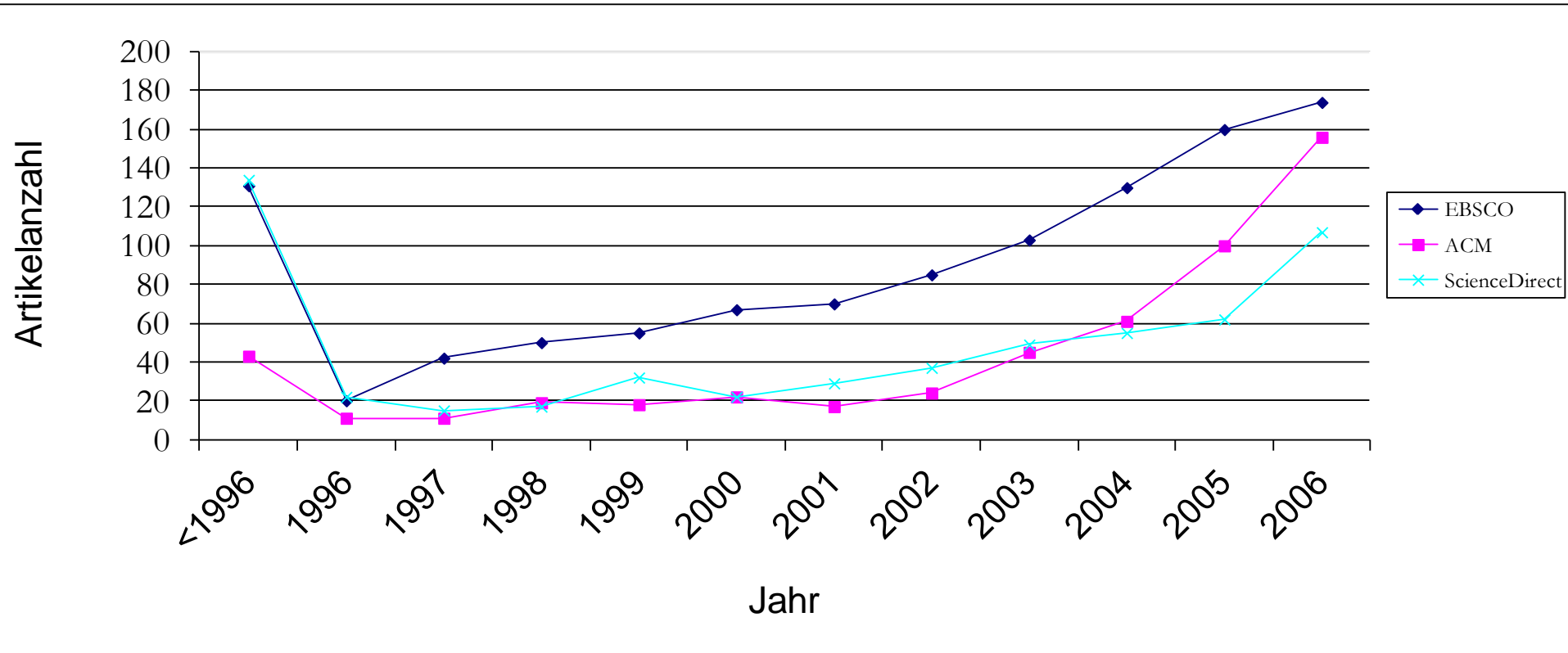
# SNA – A Recent Trend in Social Sciences Research

- Keyword search for „social“ + „network“ in 14 literature databases



*Source: Knoke, David (2007) Introduction to Social Network Analysis*

# SNA – A Recent Trend in IS Research

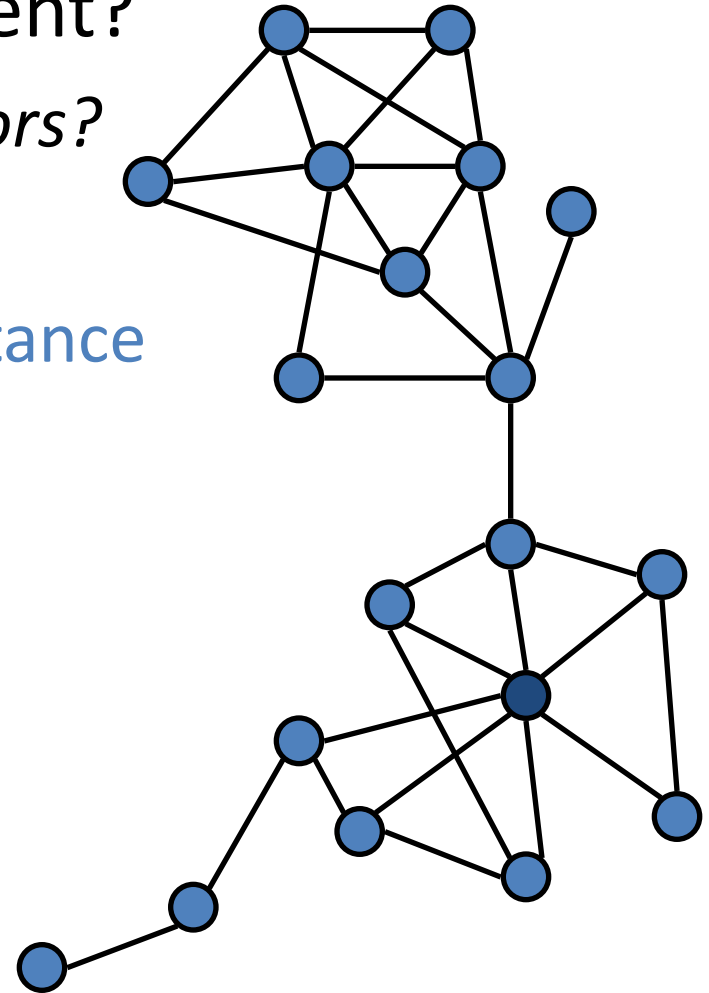




# How to analyze Social Networks?

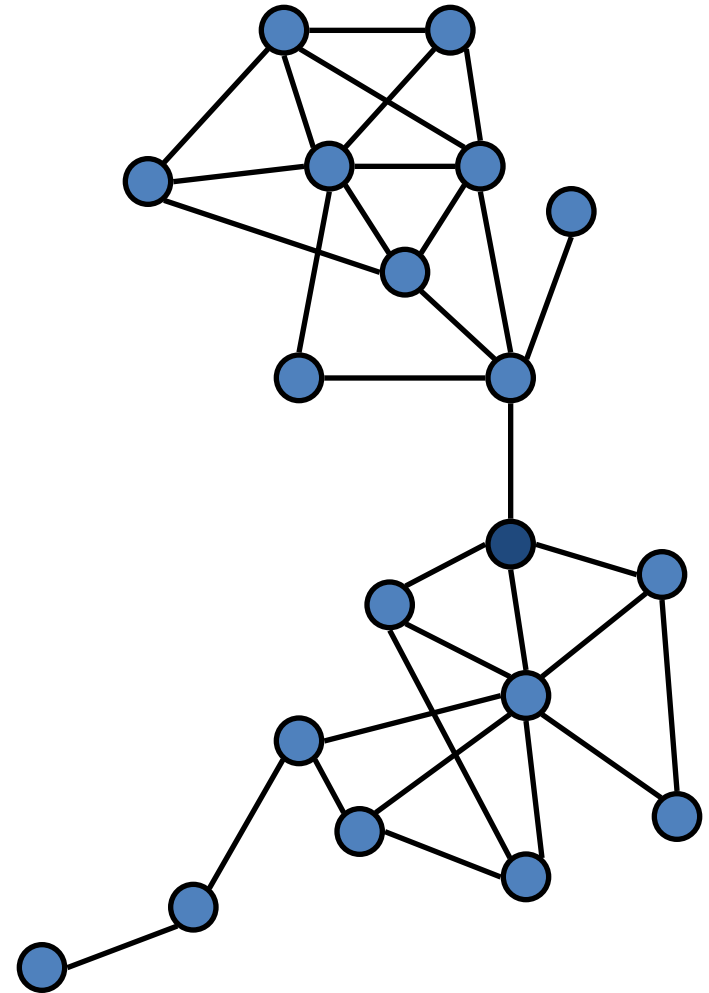
# Example: Centrality Measures

- Who is the most prominent?
  - *Who knows the most actors?*  
(Degree Centrality)
  - Who has the shortest distance to the other actors?
  - Who controls knowledge flows?
  - ...



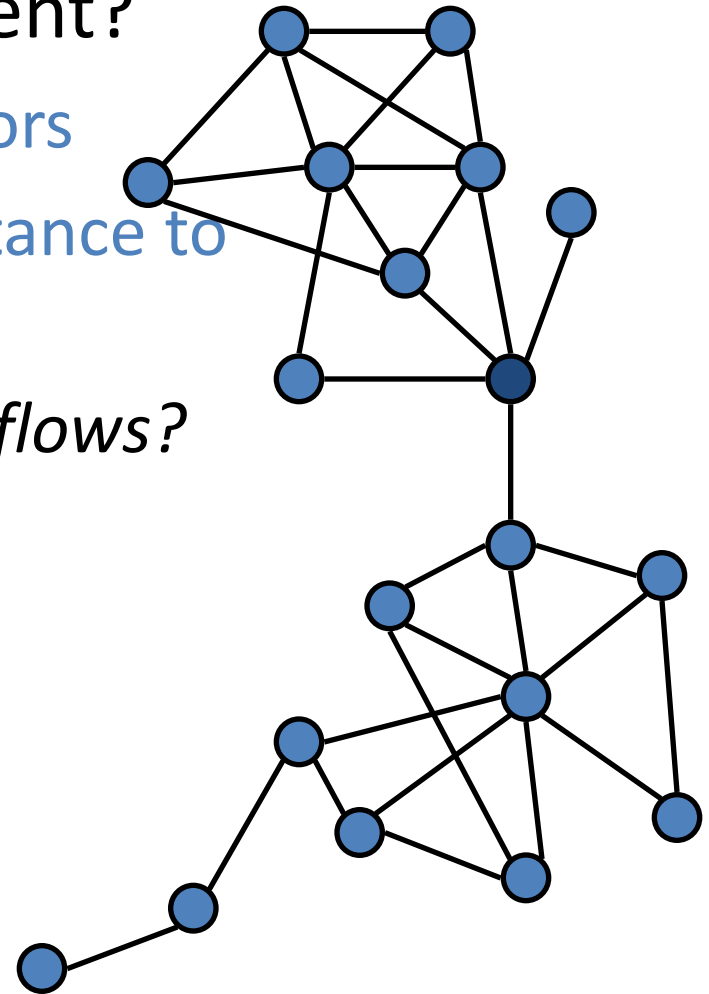
# Example: Centrality Measures

- Who is the most prominent?
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  - Who has the shortest distance to the other actors? (*Closeness Centrality*)
  - Who controls knowledge flows?
  - ...



# Example: Centrality Measures

- Who is the most prominent?
  - Who knows the most actors
  - Who has the shortest distance to the other actors?
- *Who controls knowledge flows?*  
*(Betweenness Centrality)*
- ...



# Basic Concepts

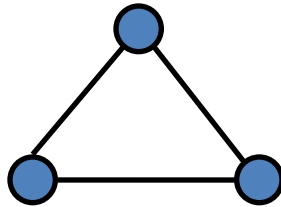
# Dyads, Triads and Relations



- *actor*



- *dyad*



- *triad*

friendship

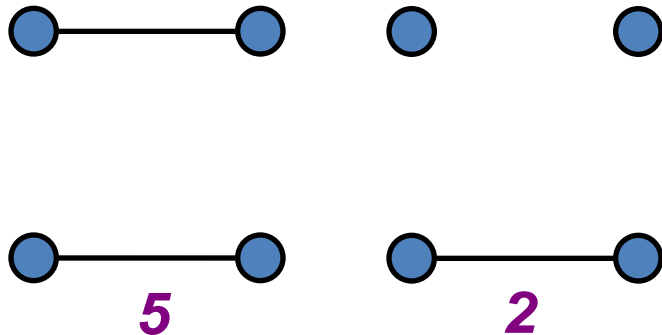
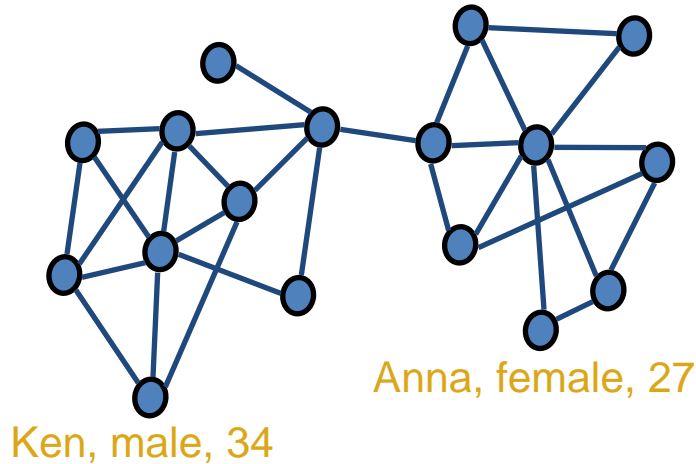


kinship



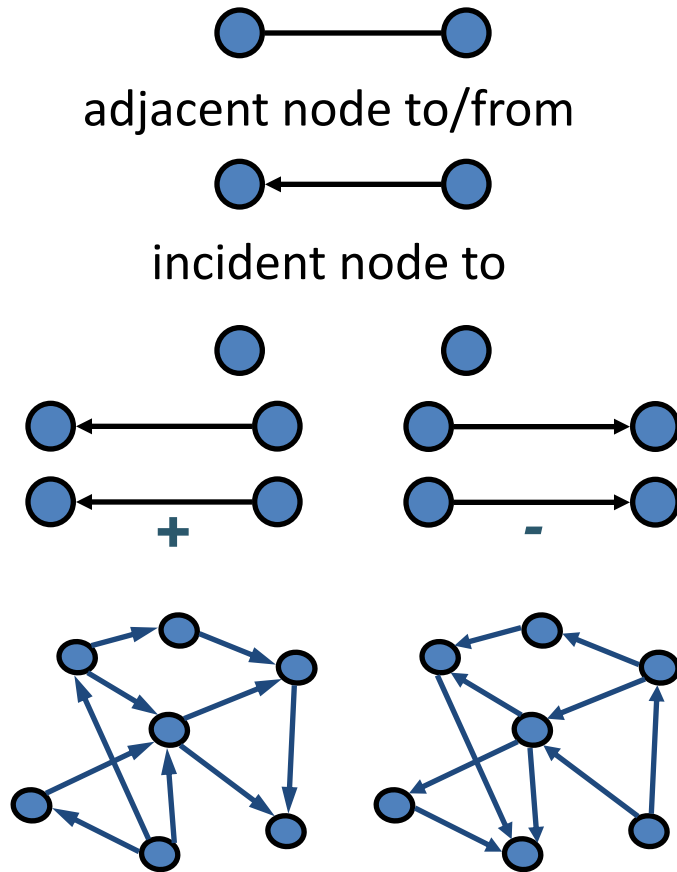
- *relation:*
  - collection of specific ties among members of a group

# Strength of a Tie



- *Social network*
  - finite set of actors and relation(s) defined on them
  - depicted in *graph/ sociogram*
    - *labeled graph*
- *Strength of a Tie*
  - *dichotomous vs. valued*
    - depicted in *valued graph* or *signed graph (+/-)*

# Strength of a Tie

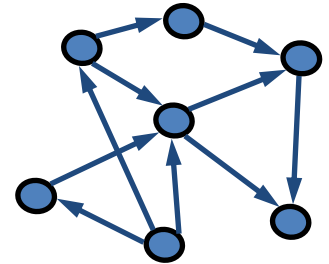


- *Strength of a Tie*
  - *nondirectional vs. directional*
    - depicted in *directed graphs (digraphs)*
    - nodes connected by *arcs*
    - *3 isomorphism classes*
      - *null dyad*
      - *mutual / reciprocal / symmetrical dyad*
      - *asymmetric / antisymmetrical dyad*
  - *converse of a digraph*
    - *reverse direction of all arcs*

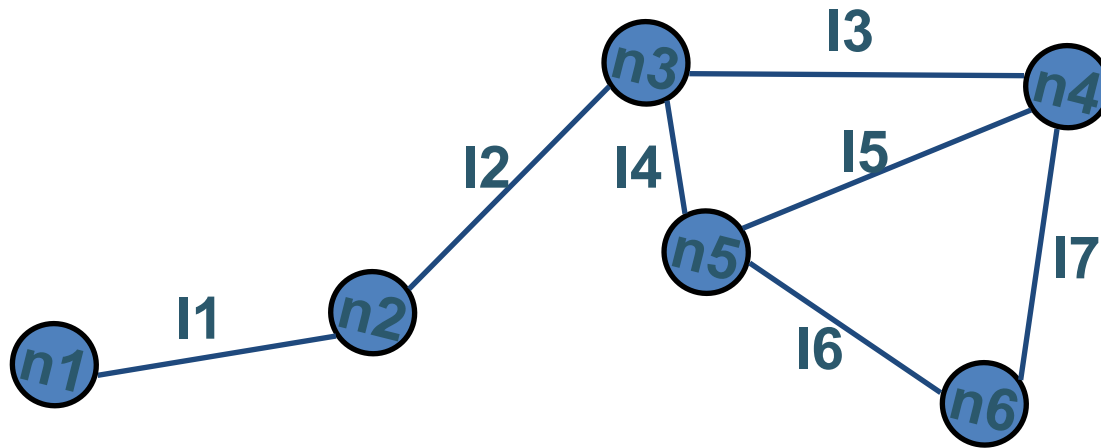


# Walks, Trails, Paths

- *(Directed) Walk (W)*
  - sequence of nodes and lines starting and ending with (different) nodes (called *origin* and *terminus*)
  - Nodes and lines can be included more than once
- *Inverse of a (directed) walk ( $W^{-1}$ )*
  - Walk in opposite order
- *Length of a walk*
  - How many lines occur in the walk? (same line counts double, in weighted graphs add line weights)
- *(Directed) Trail*
  - Is a walk in which all lines are distinct
- *(Directed) Path*
  - Walk in which all nodes and all lines are distinct
- Every path is a trail and every trail is a walk



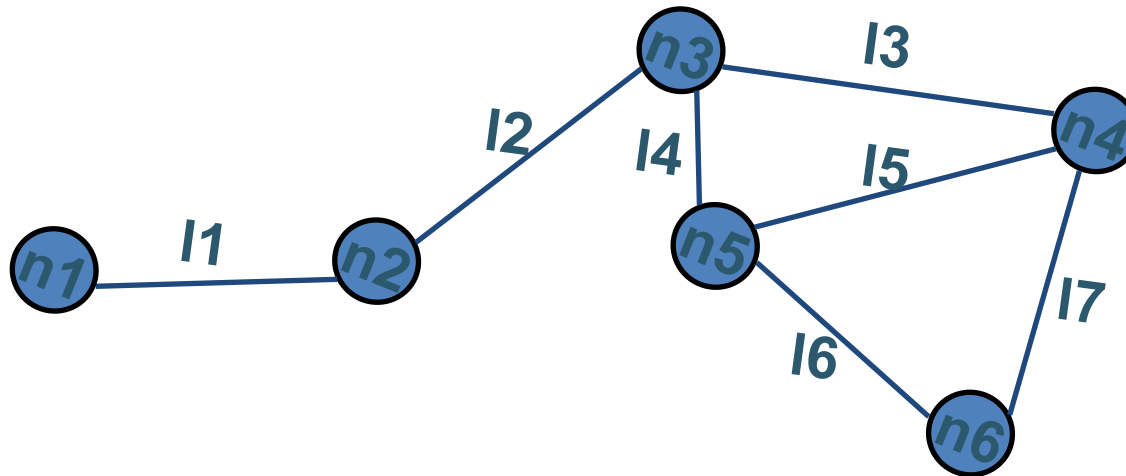
# Walks, Trails and Paths - Repetition



- $W = n1 \ l1 \ n2 \ l2 \ n3 \ l4 \ n5 \ l6 \ n6$ 
  - $n1$
  - $n3$
- $W = n1 \ l1 \ n2 \ l2 \ n3 \ l4 \ n5 \ l4 \ n3$
- $W = n1 \ l1 \ n2 \ l2 \ n3 \ l4 \ n5 \ l5 \ n4 \ l3 \ n3$
- Path
  - origin
  - terminus
- Walk
- Trail

# Reachability, Distances and Diameter

- *Reachability*
  - If there is a path between nodes  $n_i$  and  $n_j$
- *Geodesic*
  - Shortest path between two nodes
- *(Geodesic) Distance*  $d(i,j)$ 
  - Length of Geodesic (also called „degrees of separation“)



# Mathematical Notation and Fundamentals

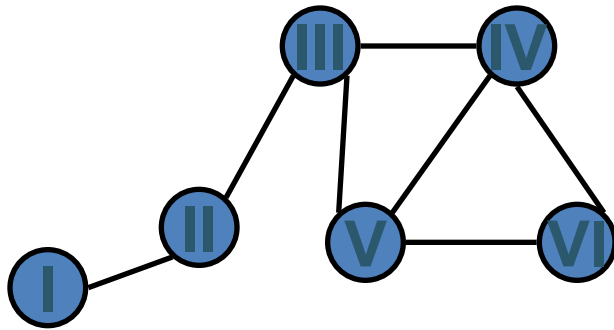
# Three different notational schemes

1. Graph theoretic
2. Sociometric
3. Algebraic

# 1. Graph Theoretic Notation

- $N$  Actors  $\{n_1, n_2, \dots, n_g\}$
- $n_i \rightarrow n_j$  there is a tie between the ordered pair  $\langle n_i, n_j \rangle$
- $n_i \not\rightarrow n_j$  there is no tie
- $(n_i, n_j)$  nondirectional relation
- $\langle n_i, n_j \rangle$  directional relation
- $g(g-1)$  number of ordered pairs in  $\langle n_i, n_j \rangle$  directional network
- $g(g-1)/2$  number of ordered pairs in nondirectional network
- $L$  collection of ordered pairs with ties  $\{l_1, l_2, \dots, l_g\}$
- $G$  graph described by sets  $(N, L)$
- *Simple graph* has no reflexive ties, loops

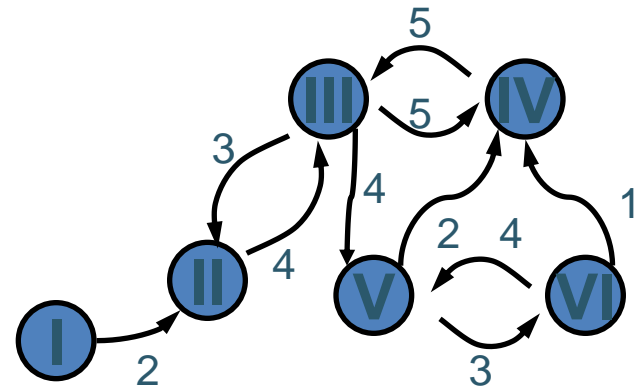
## 2. Sociometric Notation - From Graphs to (Adjacency/Socio)-Matrices



Binary, **u**ndirected

	I	II	III	IV	V	VI
I	-					
II	1	-				
III		1	-			
IV			1	-		
V			1	1	-	
VI				1	1	-

symmetrical



Valued, **d**irected

	I	II	III	IV	V	VI
I		2	0	0	0	0
II	0	0	4	0	0	0
III	0	3	0	5	4	0
IV	0	0	5	0	0	0
V	0	0	0	2	0	3
VI	0	0	0	1	4	0

## 2. Sociometric Notation

- $X$   $g \times g$  *sociomatrix* on a single relation  
 $g \times g \times R$  *super-sociomatrix* on  $R$  relations
  - $X_R$  sociomatrix on relation  $R$
- $X_{ij(r)}$  value of tie from  $n_i$  to  $n_j$  (on relation  $\chi_r$ ) where  $i \neq j$



## 2. Sociometric Notation – From Matrices to Adjacency Lists and Arc Lists

### Arc List

	I	II	III	IV	V	VI
I	-	1				
II	1	-	1			
III		1	-	1	1	
IV			1	-	1	1
V			1	1	-	1
VI				1	1	-

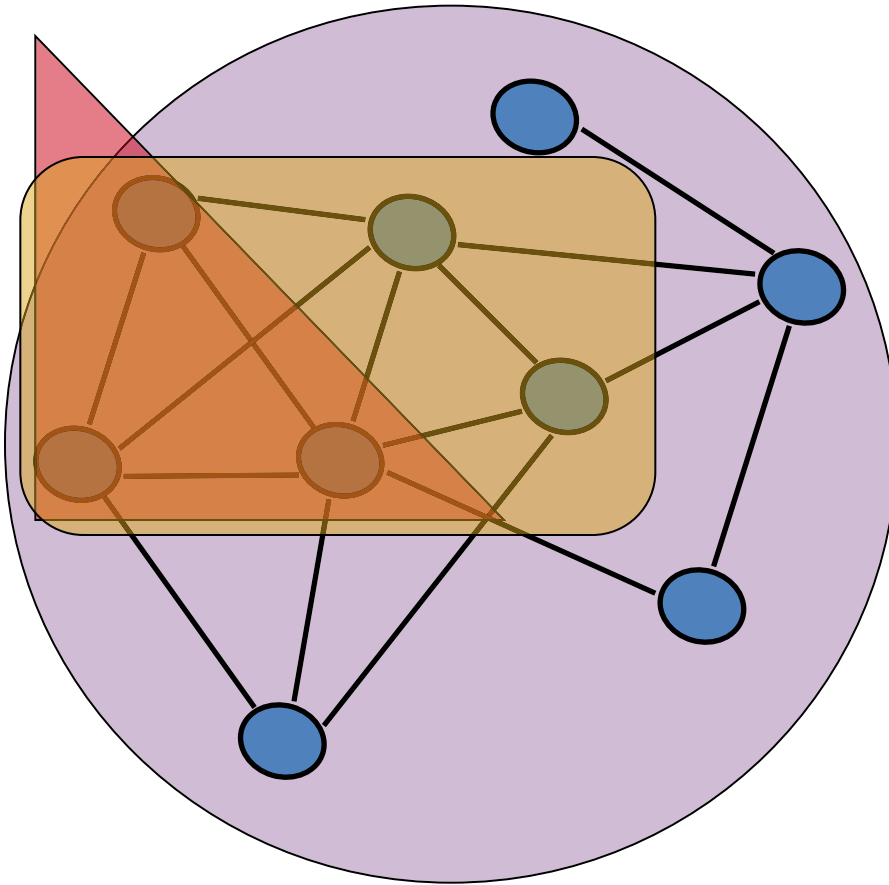
### Adjacency List

I		II
II		I III
III		II IV V
IV		III V VI
V		III V VI
VI		IV V

I II  
 II I  
 II III  
 III II  
 III IV  
 III V  
 IV III  
 IV V  
 IV VI  
 V III  
 V IV  
 V VI  
 VI IV  
 VI IV

# Network Statistics

# Different Levels of Analysis

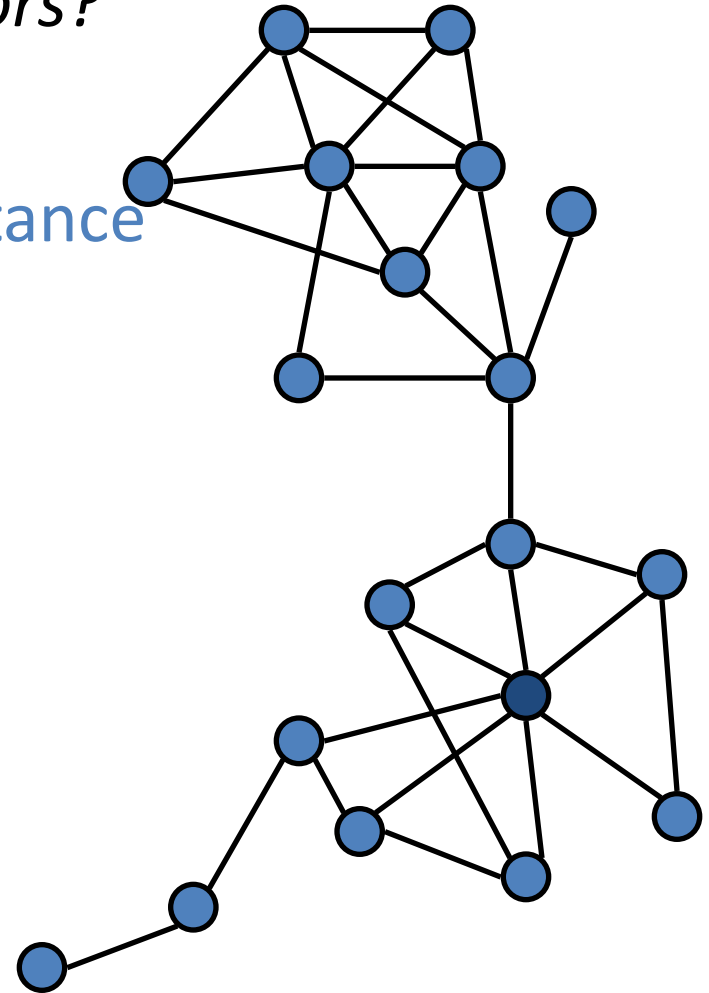


- Actor-Level
- Dyad-Level
- Triad-Level
- Subset-level (cliques / subgraphs)
- Group (i.e. global) level

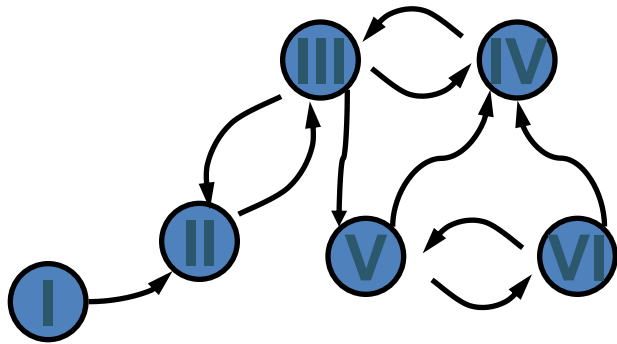
Measures at the Actor-Level:  
Measures of Prominence: Centrality and  
Prestige

# Degree Centrality

- *Who knows the most actors?*  
*(Degree Centrality)*
- Who has the shortest distance to the other actors?
- Who controls knowledge flows?
- ...



# Degree Centrality I



	I	II	III	IV	V	VI	
I		1	0	0	0	0	1
II	0	0	1	0	0	0	1
III	0	1	0	1	1	0	3
IV	0	0	1	0	0	0	1
V	0	0	0	1	0	1	2
VI	0	0	0	1	0	1	2
	0	2	2	3	1	2	

- *Indegree*  $d_I(n_i)$

- Popularity, status, deference, degree prestige

$$C_{DI}(n_i) = d_I(n_i) = \sum_j x_{ji} = x_{+i}$$

- *Outdegree*  $d_O(n_i)$

- Expansiveness

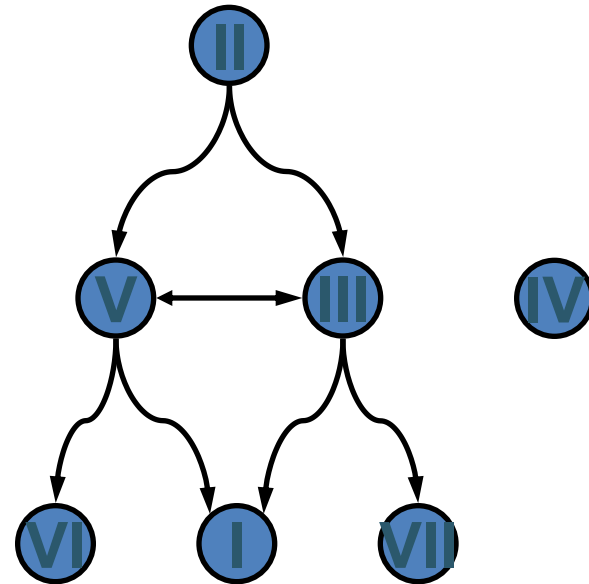
$$C_{DO}(n_i) = d_O(n_i) = \sum_j x_{ij} = x_{i+}$$

- Total degree  $\equiv 2 \times$  number of edges

**Marginals of adjacency matrix**

# Degree Centrality II

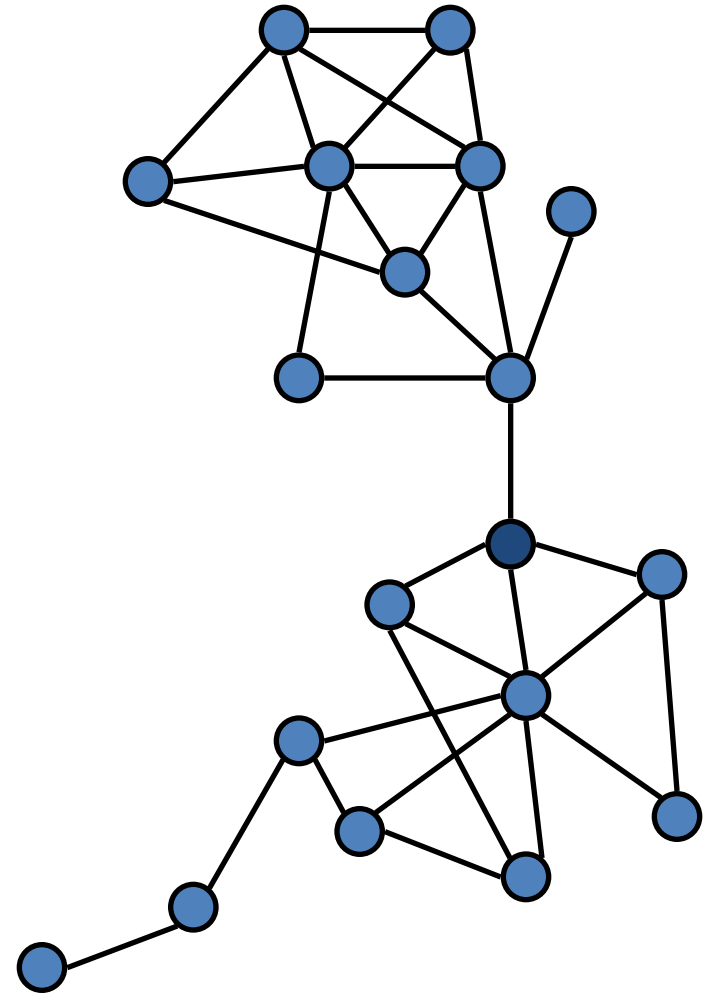
- Interpretation: opportunity to (be) influence(d)
- Classification of Nodes
  - *Isolates*
    - $d_I(n_i) = d_O(n_i) = 0$
  - *Transmitters*
    - $d_I(n_i) = 0$  and  $d_O(n_i) > 0$
  - *Receivers*
    - $d_I(n_i) > 0$  and  $d_O(n_i) = 0$
  - *Carriers / Ordinaries*
    - $d_I(n_i) > 0$  and  $d_O(n_i) > 0$
- Standardization of  $C_D$  to allow comparison across networks of different sizes: divide by its maximum value



$$C'_D(n_i) = \frac{d(n_i)}{g-1}$$

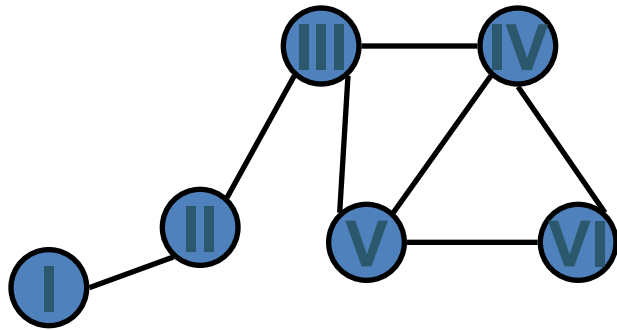
# Closeness Centrality

- Who knows the most actors?
- *Who has the shortest distance to the other actors? (Closeness Centrality)*
- Who controls knowledge flows?
- ...





# Closeness Centrality



	I	II	III	IV	V	VI	
I	-	1	2	3	3	4	13
II	1	-	1	2	2	3	9
III	2	1	-	1	1	2	7
IV	3	2	1	-	1	1	8
V	3	2	1	1	-	1	8
VI	4	3	2	1	1	-	11

- Index of expected arrival time

$$C_C(n_i) = \frac{1}{\sum_{j=1}^g d(n_i, n_j)}$$

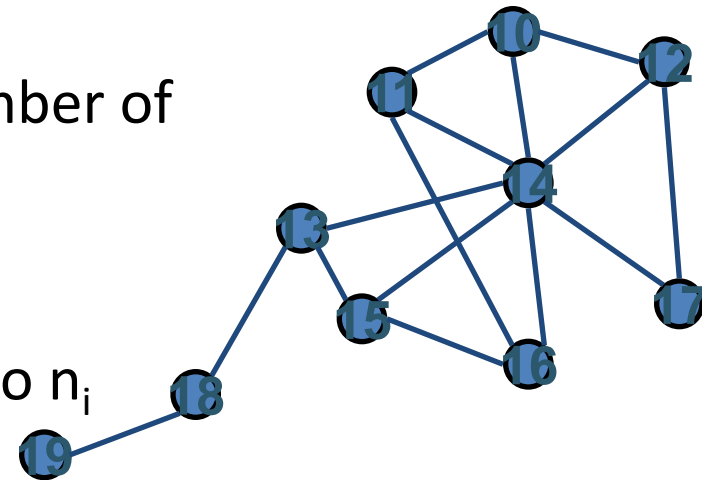
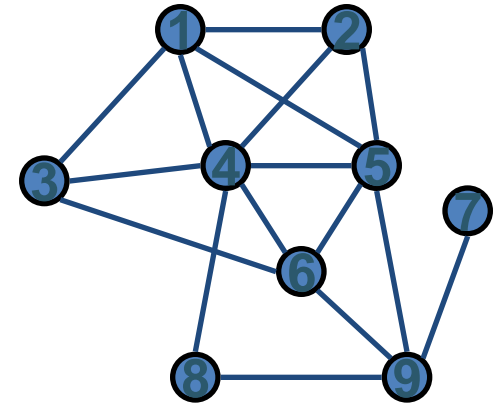
**Reciprocal of marginals of geodesic distance matrix**

- Standardize by multiplying  $(g-1)$
- Problem: Only defined for connected graphs

# Proximity Prestige

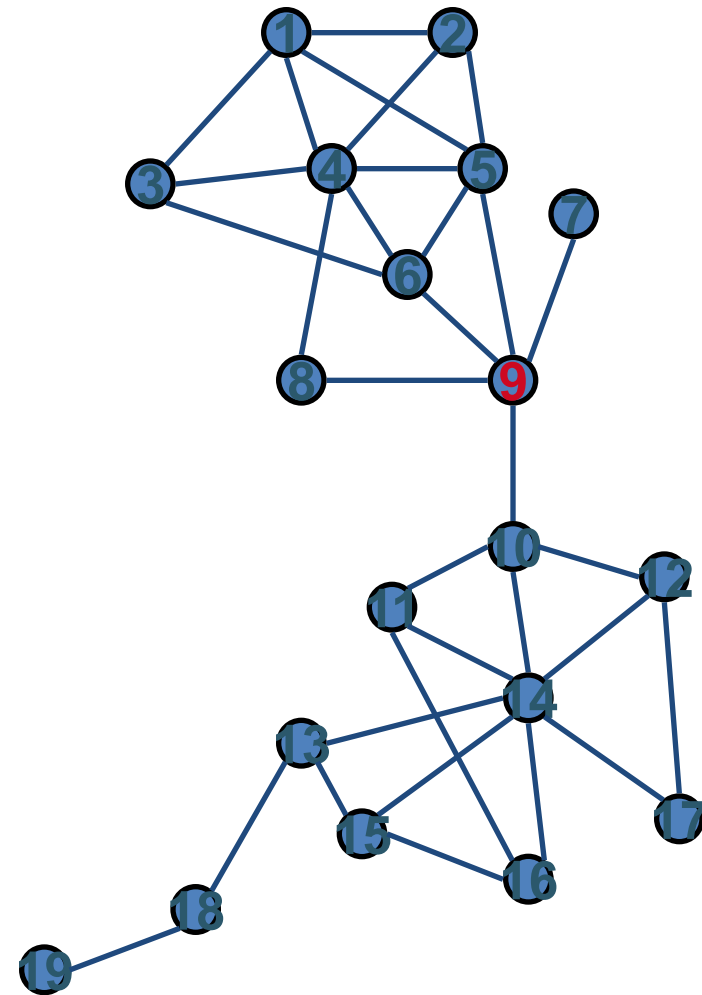
$$P_P(n_i) = \frac{I_i / (g - 1)}{\sum_{j=1}^g d(n_j, n_i) / I_i}$$

- $I_i / (g - 1)$ 
  - number of actors in the influence domain of  $n_i$
  - normed by maximum possible number of actors in influence domain
- $\sum d(n_j, n_i) / I_i$ 
  - average distance these actors are to  $n_i$



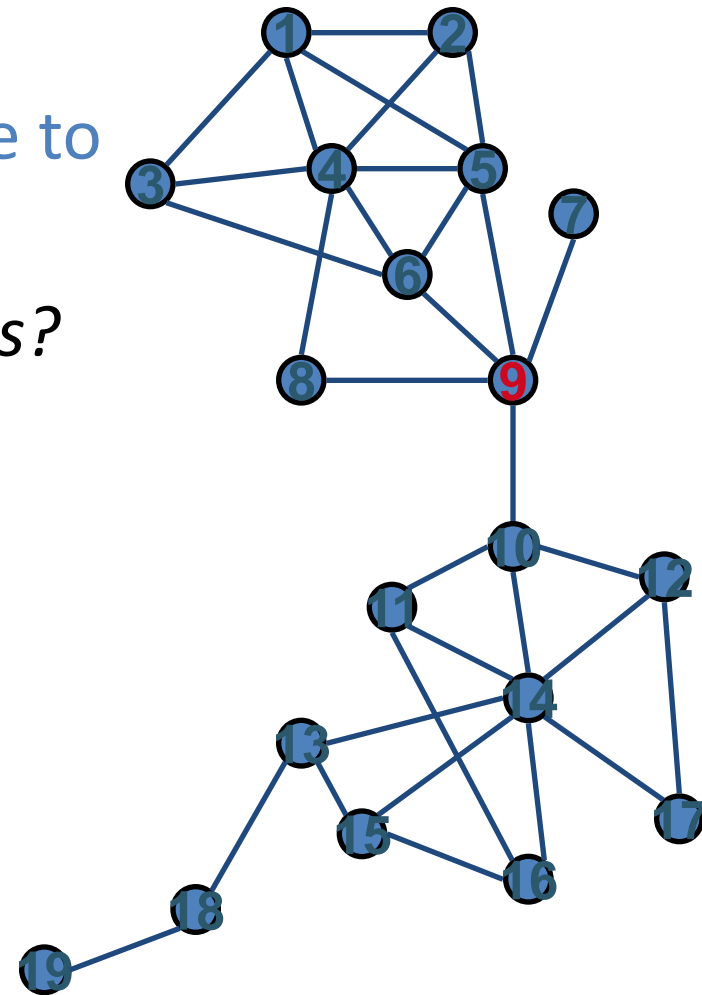
# Eccentricity / Association Number

- Largest geodesic distance between a node and any other node
- $\max_j d(i,j)$



# Betweenness Centrality

- Who knows the most actors?
- Who has the shortest distance to the other actors?
- *Who controls knowledge flows?  
(Betweenness Centrality)*
- ...

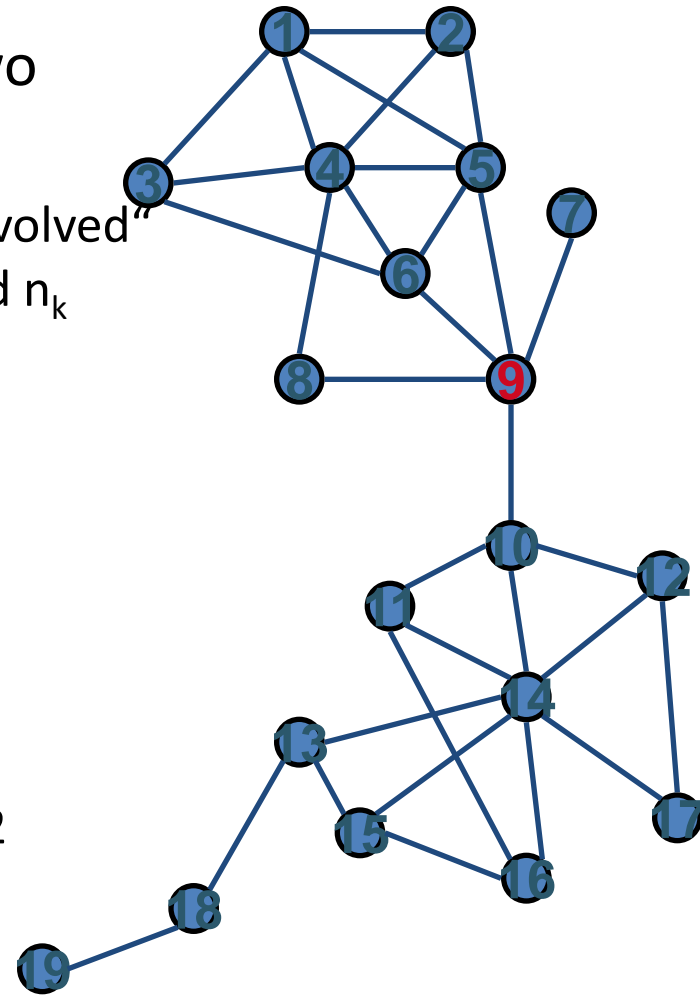


# Betweenness Centrality

- How many geodesic linkings between two actors  $j$  and  $k$  contain actor  $i$ ?
  - $g_{jk}(n_i)/g_{jk}$  probability that distinct actor  $n_i$  „involved“ in communication between two actors  $n_j$  and  $n_k$

$$C_B(n_i) = \frac{\sum_{j < k} g_{jk}(n_i)}{g_{jk}}$$

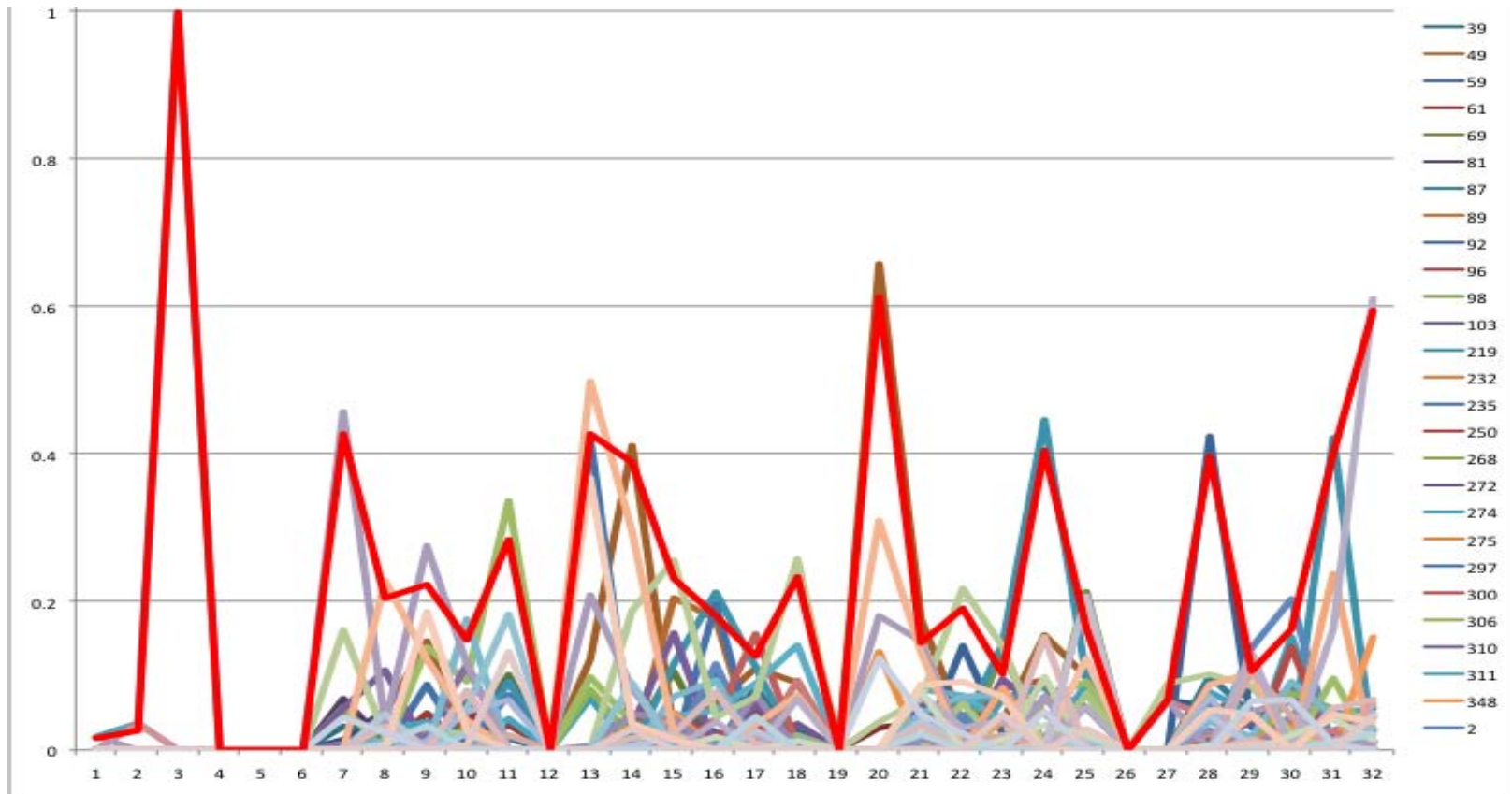
- standardized by dividing through  $(g-1)(g-2)/2$



# Several other Centrality Measures

- ...beyond the scope of this lecture
  - *Status or Rank Prestige, Eigenvector Centrality*
    - also reflects status or prestige of people whom actor is linked to
    - Appropriate to identify *hubs* (actors adjacent to many peripheral nodes) and *bridges* (actors adjacent to few central actors)
      - attention: more common, different meaning of bridge!!!
  - *Information Centrality*
  - see Wasserman & Faust (1994), p. 192 ff.
  - *Random Walk Centrality*
    - see Newman (2005)

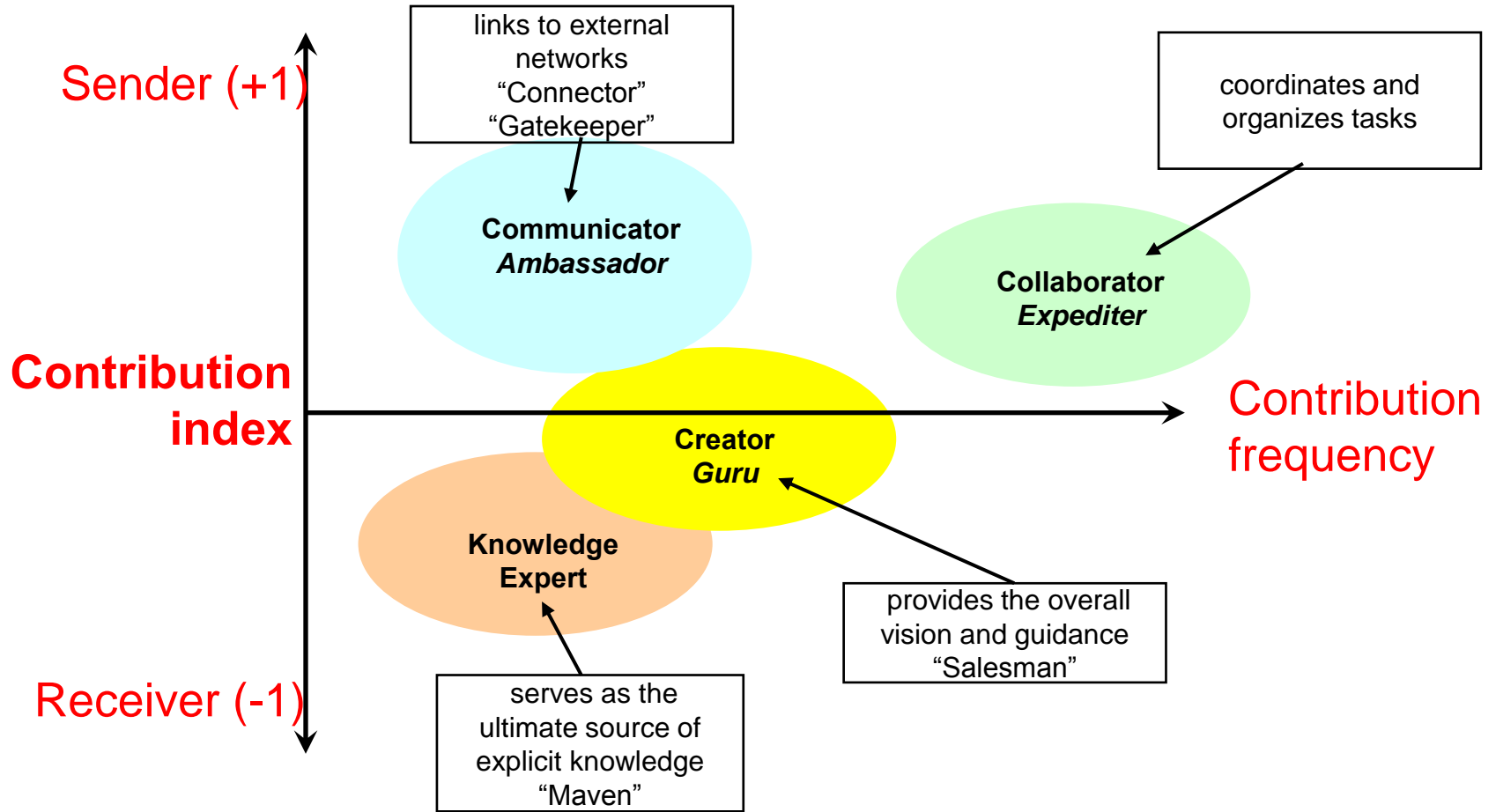
# Condor – Betweenness Centrality



# (Actor) Contribution Index

*messa g essen t - messa g esreceived*

*messa g essen t + messa g esreceived*

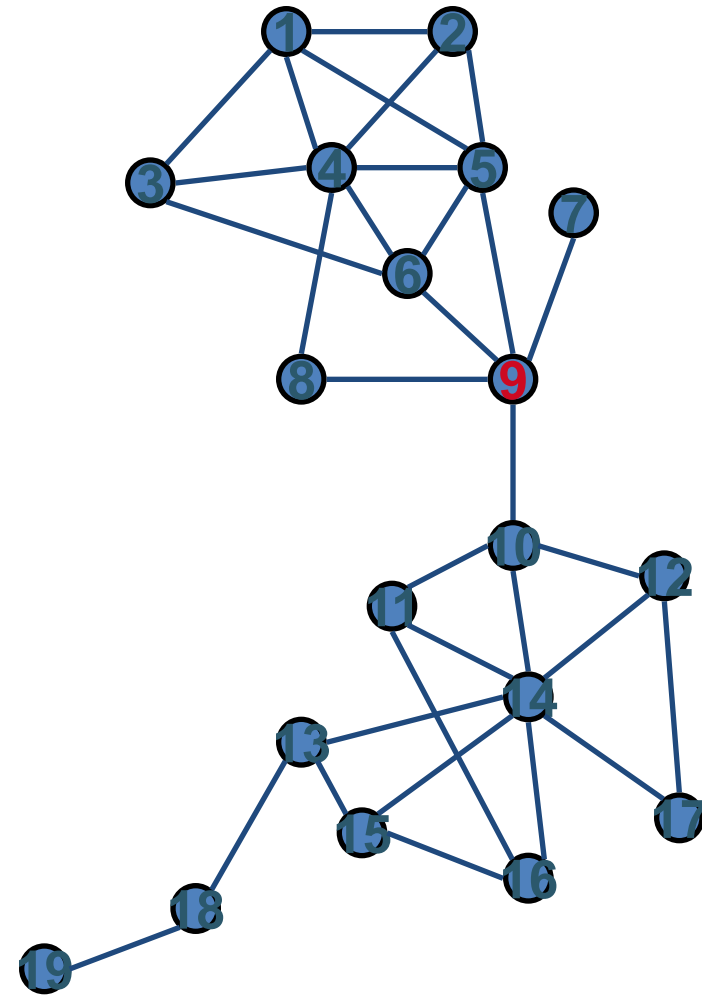




# Measures at the Group-(Global-)Level and Subgroup-Level

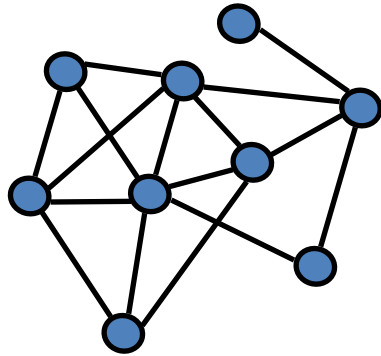
# Diameter of a Graph and Average Geodesic Distance

- Diameter
  - Largest geodesic distance between any pair of nodes
- Average Geodesic Distance
  - How fast can information get transmitted?

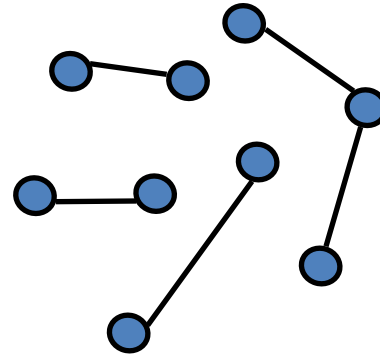


# Density

- Proportion of ties in a graph

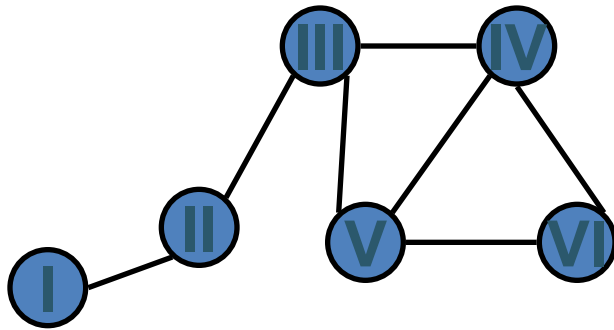


High density (44%)



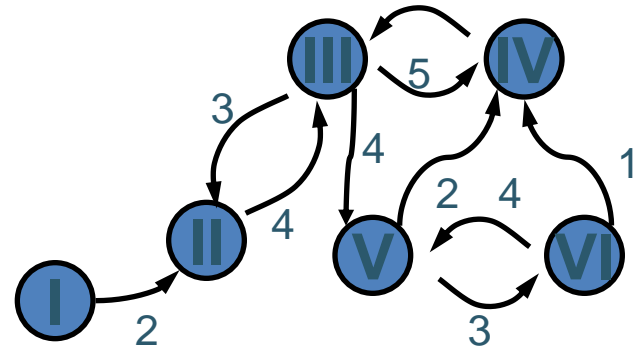
Low density (14%)

# Density



$$\Delta = \frac{L}{g(g-1)/2} = \frac{L}{\binom{g}{2}}$$

In undirected graph:  
Proportion of ties



$$\Delta = \frac{\sum_{i=1}^g \sum_{j=1}^g x_{ij}}{g(g-1)}$$

In valued directed graph:  
Average strength of the arcs

# Group Centralization I

- How equal are the individual actors' centrality values?
  - $C_A(n_i^*)$  actor centrality index
  - $C_A(n^*)$   $\max_i C_A(n_i^*)$
  - $\sum_{i=1}^g [C_A(n^*) - C_A(n_i)]$  sum of difference between largest value and observed values
- General centralization index:

$$C_A = \frac{\sum_{i=1}^g [C_A(n^*) - C_A(n_i)]}{\max \sum_{i=1}^g [C_A(n^*) - C_A(n_i)]}$$

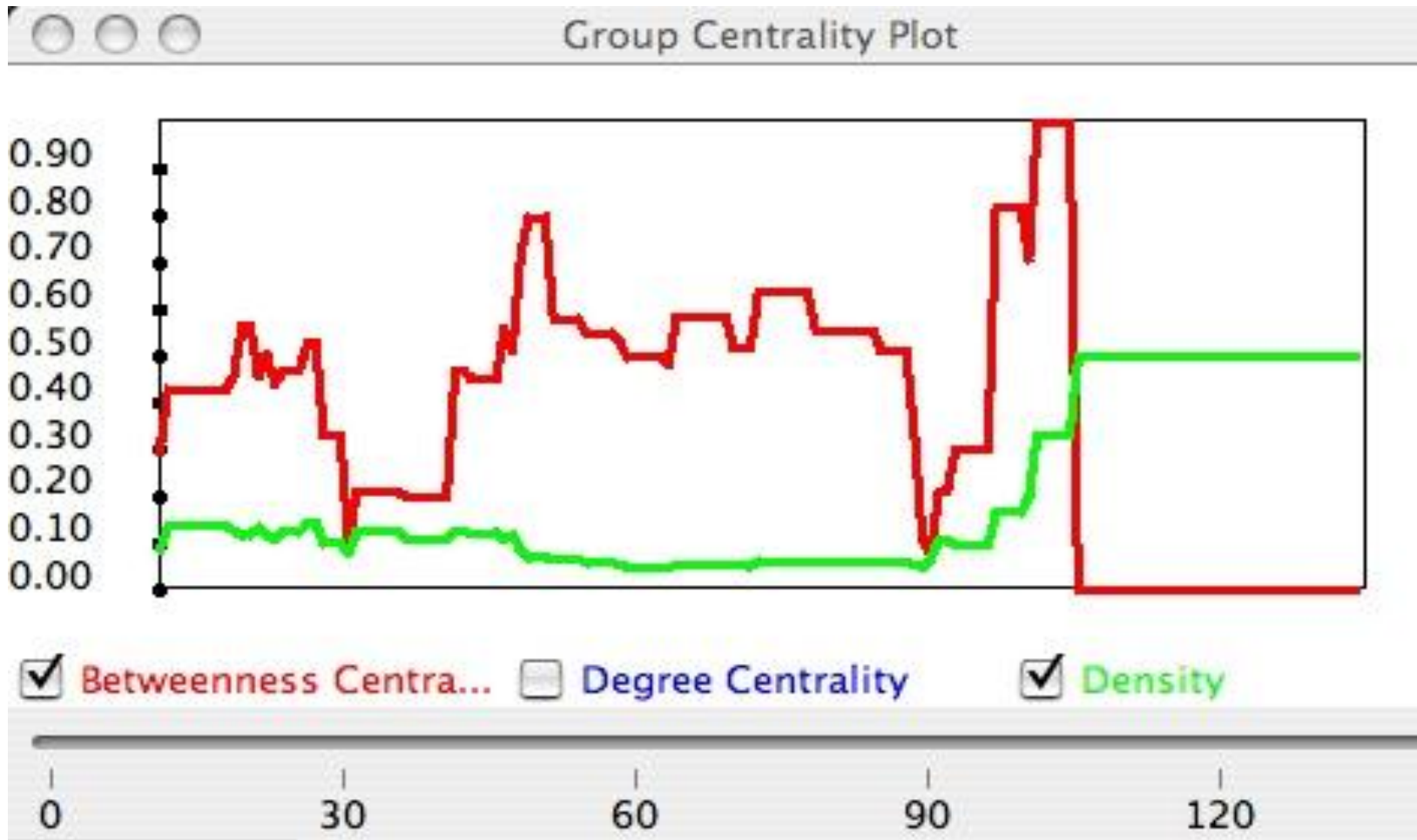
# Group Centralization II

$$C_D = \frac{\sum_{i=1}^g [C_D(n^*) - C_D(n_i)]}{(g-1)(g-2)}$$

$$C_C = \frac{\sum_{i=1}^g [C'_C(n^*) - C'_C(n_i)]}{[(g-1)(g-2)](2g-3)}$$

$$C_B = \frac{\sum_{i=1}^g [C_B(n^*) - C_B(n_i)]}{(g-1)^2(g-2)} = \frac{\sum_{i=1}^g [C'_B(n^*) - C'_B(n_i)]}{(g-1)}$$

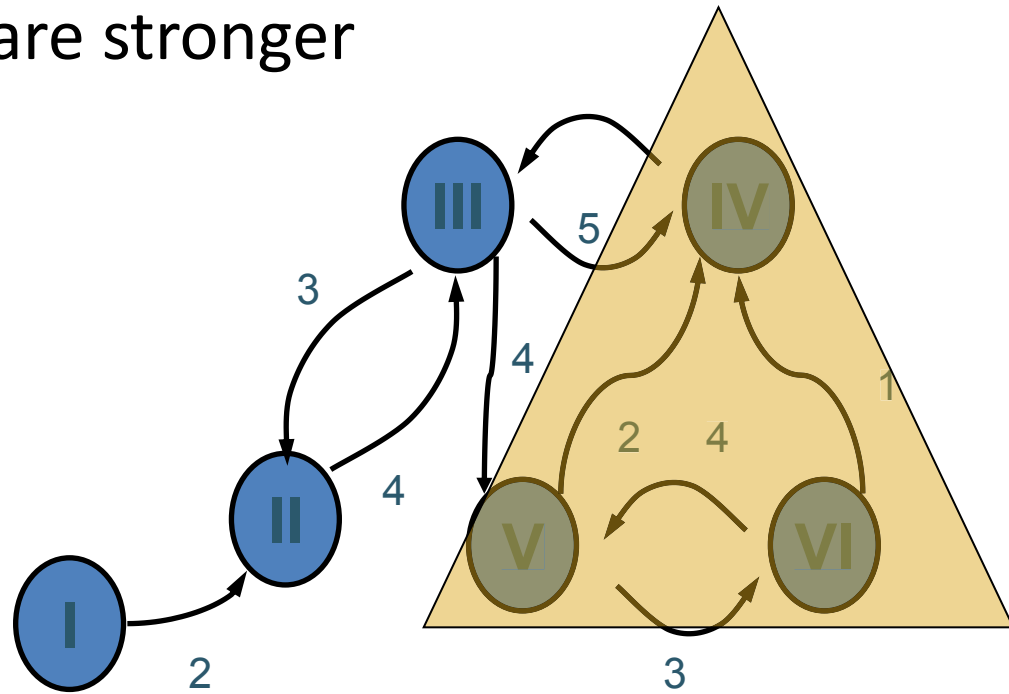
# Condor – Group Centralization



# Subgroup Cohesion

- average strength of ties within the subgroup divided by average strength of ties that are from subgroup members to outsiders
- $>1 \rightarrow$  ties in subgroup are stronger

$$\frac{\sum_{i \in N_s} \sum_{j \in N_s} x_{ij}}{g_s (g_s - 1)}$$
$$\frac{\sum_{i \in N_s} \sum_{j \notin N_s} x_{ij}}{g_s (g - g_s)}$$



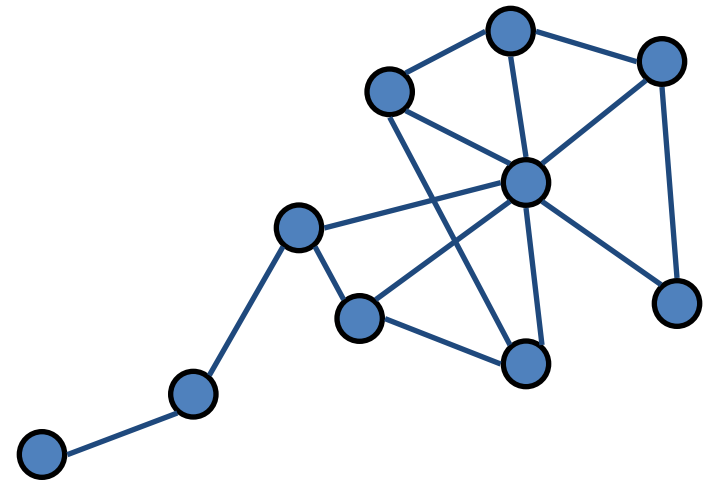
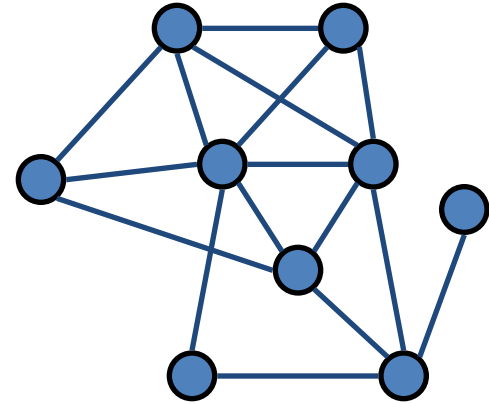


# Connectivity of Graphs and Cohesive Subgroups

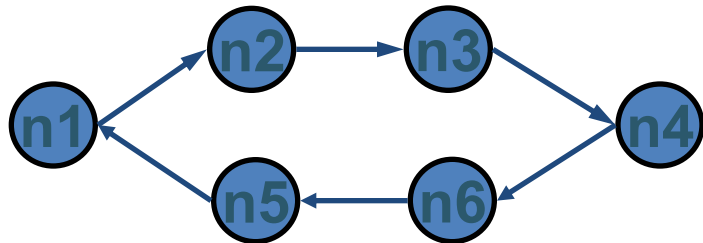
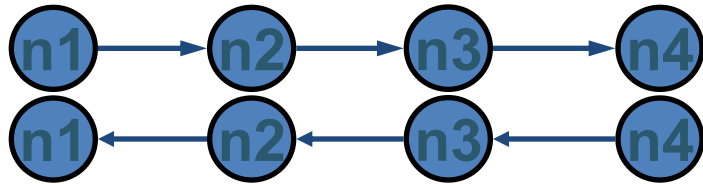
# Connectivity of Graphs

# Connected Graphs, Components, Cutpoints and Bridges

- *Connectedness*
  - A graph is connected if there is a path between every pair of nodes
- *Components*
  - Connected subgraphs in a graph
  - Connected graph has 1 component
  - Two disconnected graphs are one social network!!!



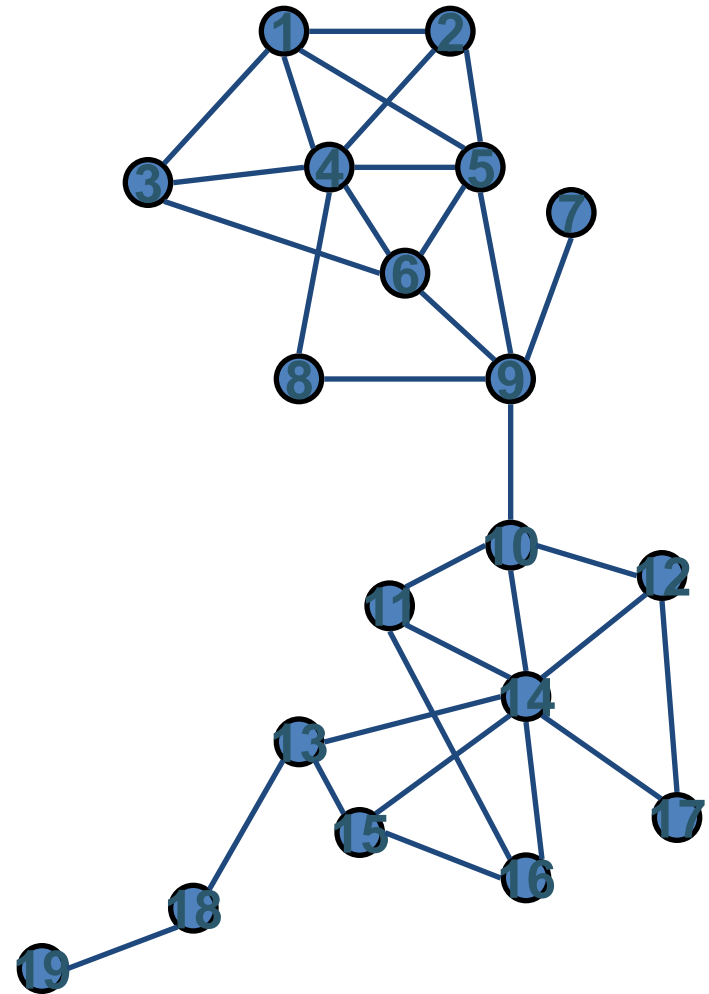
# Connected Graphs, Components, Cutpoints and Bridges



- Connectivity of pairs of nodes and graphs
  - *Weakly connected*
    - Joined by semipath
  - *Unilaterally connected*
    - Path from  $n_j$  to  $n_i$  or from  $n_i$  to  $n_j$
  - *Strongly connected*
    - Path from  $n_j$  to  $n_i$  *and* from  $n_i$  to  $n_j$
    - Path may contain different nodes
  - *Recursively Connected*
    - Nodes are strongly connected and both paths use the same nodes and arcs in reverse order

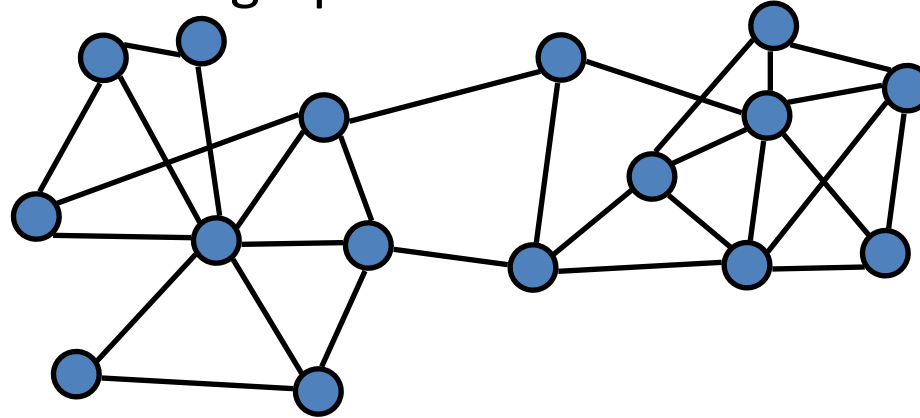
# Connected Graphs, Components, Cutpoints and Bridges

- *Cutpoints*
  - number of components in the graph that contain node  $n_j$  is fewer than number of components in subgraphs that results from deleting  $n_j$  from the graph
- *Cutsets (of size  $k$ )*
  - $k$ -node cut
- *Bridges / line cuts*
  - Number of components...that contain line  $l_k$



# Node- and Line Connectivity

- How vulnerable is a graph to removal of nodes or lines?



*Point connectivity /  
Node connectivity*

- Minimum number of  $k$  for which the graph has a  $k$ -node cut
- For any value  $<k$  the graph is  $k$ -node-connected

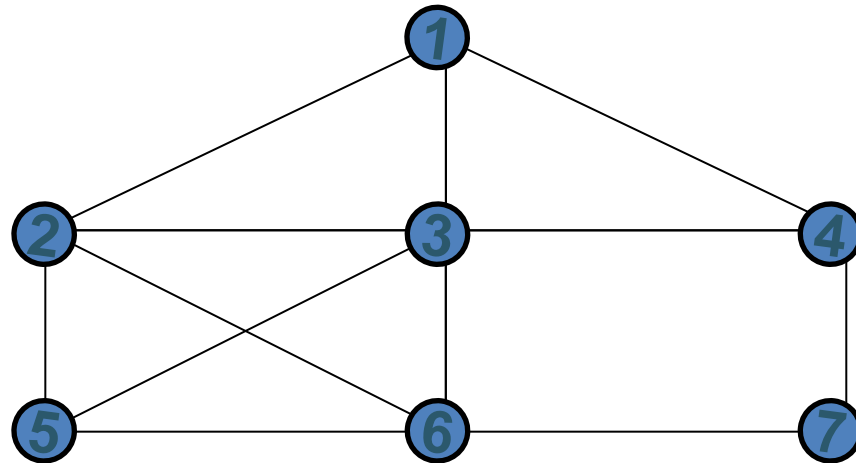
*Line connectivity / Edge  
connectivity*

- Minimum number  $\lambda$  for which graph has a  $\lambda$ -line cut

# Cohesive Subgroups

# Cohesive Subgroups, (n-)Cliques, n-Clans, n-Clubs, k-Plexes, k-Cores

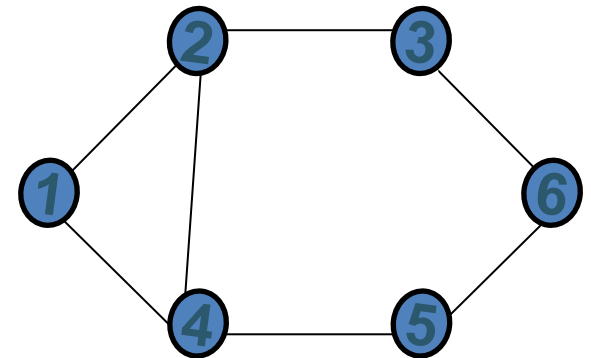
- *Cohesive Subgroup*
  - Subset of actors among there are relatively strong, direct, intense, frequent or positive ties
- *Complete Graph*
  - All nodes are adjacent
- *Clique*
  - Maximal complete subgraph of three or more nodes
  - Cliques can overlap
  - {1, 2, 3}
  - {1, 3, 4}
  - {2, 3, 5, 6}





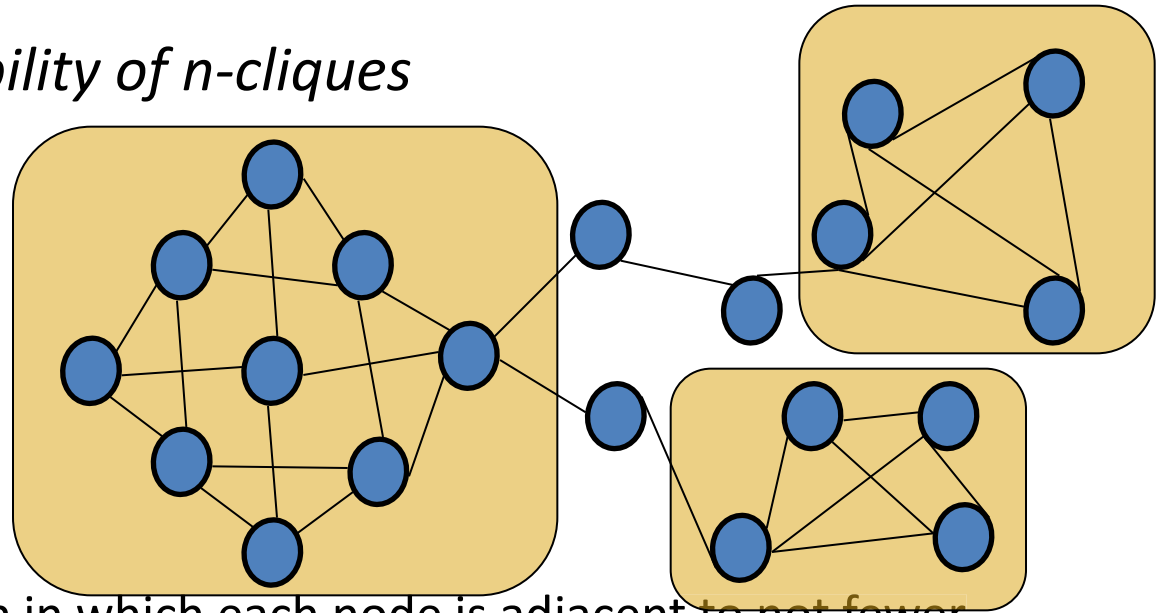
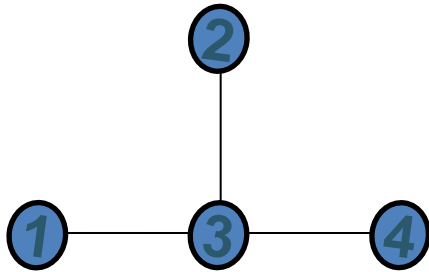
# Cohesive Subgroups, (n-)Cliques, n-Clans, n-Clubs, k-Plexes, k-Cores

- *n-clique*
  - maximal subgraph in which  $d(i,j) \leq n$  for all  $n_i, n_j$
  - 2: cliques: {2, 3, 4, 5, 6} and {1, 2, 3, 4, 5}
  - intermediaries in geodesics do not have to be n-clique members themselves!
- *n-clan*
  - *n-clique* in which the  $d(i,j) \leq n$  for the subgraph of all nodes in the n-clique
  - 2-clan: {2, 3, 4, 5, 6}
- *n-club*
  - maximal subgraph of diameter n
  - 2-clubs: {1, 2, 3, 4}; {1, 2, 3, 5} and {2, 3, 4, 5, 6}



# Cohesive Subgroups, (n-)Cliques, n-Clans, n-Clubs, k-Plexes, k-Cores

- *Problem: vulnerability of n-cliques*



- *k-plexes*

- maximal subgraph in which each node is adjacent to not fewer than  $g_s - k$  nodes („maximal“: no other nodes in subgraph that also have  $d_s(i) \geq (g_s - k)$  ]

- *k-cores*

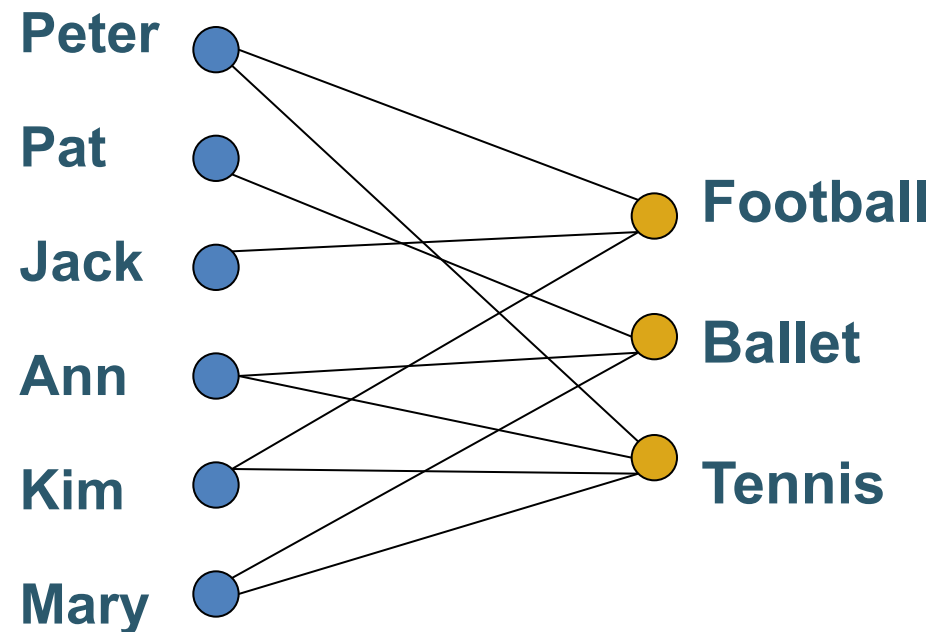
- subgraph in which each node is adjacent to at least  $k$  other nodes in the subgraph

# Analyzing Affiliation Networks

# Affiliation Matrix, Bipartite Graph and Hypergraph, Rate of Participation, Size of Events

*Two-mode network / affiliation network / membership network / hypernetwork*

- nodes can be partitioned in two subsets
  - N (for example g persons)
  - M (for example h clubs)
- depicted in *Bipartite Graph*
- lines between nodes belonging to different subsets



# Affiliation Matrix, Bipartite Graph and Hypergraph, Rate of Participation, Size of Events

## Affiliation Matrix (Incidence Matrix)

- Connections among members of one of the modes based on linkages established through second mode
- $g$  actors,  $h$  events
- $A = \{a_{ij}\}$  ( $g \times h$ )

		Event				
		Football	Ballet	Tennis	rate of participation	
Actor	Peter	1		1	1	
	Pat		1		1	
	Jack	1			1	
	Ann		1	1	2	
	Kim	1		1	2	
	Mary		1	1	2	
size of event		3	3	4		

# Affiliation Matrix, Bipartite Graph and Hypergraph, Rate of Participation, Size of Events

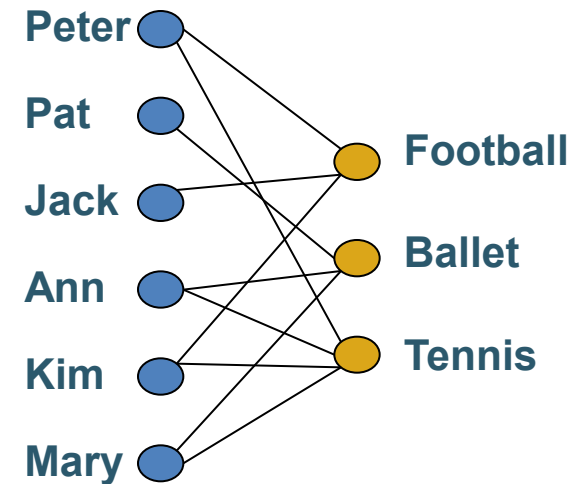
## ■ Sociomatrix [ (g+h) × (g+h) ]

	Peter	Pat	Jack	Ann	Kim	Mary	Football	Ballet	Tennis
Peter	-	0	0	0	0	0	1	0	1
Pat	0	-	0	0	0	0	0	1	0
Jack	0	0	-	0	0	0	1	0	0
Ann	0	0	0	-	0	0	0	1	1
Kim	0	0	0	0	-	0	1	0	1
Mary	0	0	0	0	0	-	0	1	1
Football	1	0	1	0	1	0	-	0	0
Ballet	0	1	0	1	0	1	0	-	0
Tennis	1	0	0	1	1	1	0	0	-

# Affiliation Matrix, Bipartite Graph and Hypergraph, Rate of Participation, Size of Events

- *Homogenous* pairs and *heterogenous* pairs
- $X_r^N (g \times g)$ ,  $X_r^M (h \times h)$ ,  $X_r^{N,M} (g \times h)$ ,  $X_r^{N,M} (h \times g)$ 
  - One-mode sociomatrices  $X^N$  [and  $X^M$ ]
    - rows, columns: actors [events];
    - $x_{ij}$ : *co-membership* [number of actors in both events] (main diagonal meaningful, e.g. total events attended by an actor)

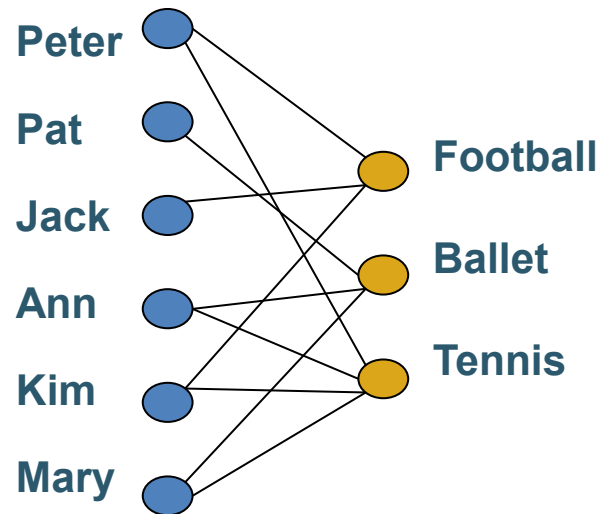
	Peter	Pat	Jack	Ann	Kim	Mary
Peter	2	0	1	1	2	1
Pat	0	1	0	1	0	1
Jack	1	0	1	0	1	0
Ann	1	0	0	2	1	1
Kim	2	0	1	1	2	1
Mary	1	0	0	1	1	2



# Affiliation Matrix, Bipartite Graph and Hypergraph, Rate of Participation, Size of Events

## ■ Event Overlap / Interlocking Matrix

	Football	Ballet	Tennis
Football	3	0	2
Ballet	0	3	1
Tennis	2	1	4





# Cohesive Subsets of Actors or Events

- *clique at level  $c$  (cf. also  $k$ -plexes,  $n$ -cliques etc.)*
  - subgraph in which all pairs of events share at least  $c$  members
- *connected at level  $q$* 
  - subset in which all actors in the path are co-members of at least  $q+1$  events

...HOLDOUT I...

# When is Which Centrality Measure Appropriate?

*Source: Borgatti, Stephen P. (2005) Centrality and Network Flow, Social Networks 27, p. 55-71*

# Assumptions of Centrality Measures

- Which things flow through a network and how do they flow?

	<b>Transfer</b>	<b>Serial</b>	<b>Parallel</b>
<b>Walks</b>	<b>Money exchange</b>	<b>Emotional support</b>	<b>Attitude influencing</b>
<b>Trails</b>	<b>Used Book</b>	<b>Gossip</b>	<b>E-mail broadcast</b>
<b>Paths</b>	<b>Mooch</b>	<b>Viral infection</b>	<b>Internet name-server</b>
<b>Geodesics</b>	<b>Package Delivery</b>	<b>Mitotic reproduction</b>	<b>&lt;no process&gt;</b>

*Source: Borgatti, Stephen P. (2005) Centrality and Network Flow, Social Networks 27, p. 55-71*

# Assumptions of Centrality Measures

- Example: Betweenness Centrality
  - Information travels along the shortest route

	Transfer	Serial	Parallel
Walks	Random Walk Betweenness	?	Closeness Degree Eigenvector
Trails	?	?	Closeness Degree
Paths	?	?	Closeness Degree
Geodesics	Closeness Betweenness	Closeness	?

# Adequacy of Centrality Measures

	Transfer	Serial	Parallel
Walks	Money exchange	Emotional support	Attitude influencing
Trails	Used Book	Gossip	E-mail broadcast
Paths	Mooch	Viral infection	Internet name-server
Geodesics	Package Delivery	Mitotic reproduction	<no process>

*Source: Borgatti, Stephen P. (2005) Centrality and Network Flow, Social Networks 27, p. 55-71*

# How to Calculate Geodesic Distance Matrices?

# From Adjacency Matrices to (Geodesic) Distance Matrices I – (Reachability)

*Repetition: Matrix Multiplication*

- $XY = Z$

**X**

3	0	2
1	4	2

**g × h**

**Y**

2	3
1	4
4	2

**h × k**

**Z**

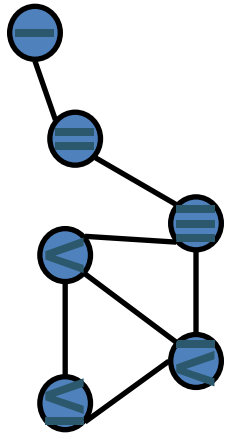
14	13
14	23

**g × k**

- $$z_{ij} = \sum_{n=1}^h x_{in} y_{nj}$$



# From Adjacency Matrices to (Geodesic) Distance Matrices II – (Reachability)



	I	II	III	IV	V	VI
I	0	1	0	0	0	0
II	1	0	1	0	0	0
III	0	1	0	1	1	0
IV	0	0	1	0	1	1
V	0	0	1	1	0	1
VI	0	0	0	1	1	0

X

	I	II	III	IV	V	VI
I	0	1	0	0	0	0
II	1	0	1	0	0	0
III	0	1	0	1	1	0
IV	0	0	1	0	1	1
V	0	0	1	1	0	1
VI	0	0	0	1	1	0

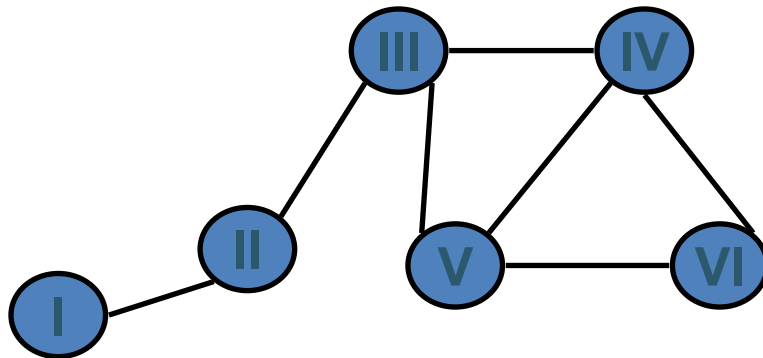
Power Matrix: Multiplying adjacency matrices

- $x_{ik}x_{kj} = 1$  only if lines  $(n_i, n_k)$  and  $(n_k, n_j)$  are present, i.e.  $X^{[2,3,4]}$  counts the number of walks  $(n_i, n_k, n_j)$  of length 1 [2,3,4] between nodes  $n_i$  and  $n_j$

	I	II	III	IV	V	VI
I	1	0	1	0	0	0
II	0	2	0	1	1	0
III	1	0	3	1	1	2
IV	0	1	1	3	2	1
V	0	1	1	2	3	1
VI	0	0	2	1	1	2

# From Adjacency Matrices to (Geodesic) Distance Matrices II – (Reachability)

- $x_{ij} > 0$  ?  
 → two nodes can be connected by paths of length  $\leq (g-1)$
- Calculate  $X^{[\Sigma]} = X + X^2 + X^3 + \dots + X^{g-1}$
- $X^{[\Sigma]}$  shows total number of walks from  $n_i$  to  $n_j$

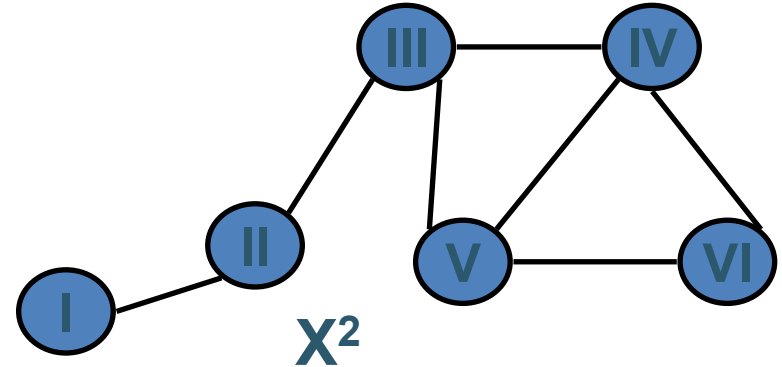


$X^2$

	I	II	III	IV	V	VI
I	1	0	1	0	0	0
II	0	2	0	1	1	0
III	1	0	3	1	1	2
IV	0	1	1	3	2	1
V	0	1	1	2	3	1
VI	0	0	2	1	1	2

# From Graphs to (Geodesic Distance)-Matrices (Reachability) – Geodesic Distance

- observer power matrices
- first power  $p$  for which the  $(i,j)$  element is non-zero gives the shortest path
- $d(i,j) = \min_p x_{ij}^{[p]} > 0$



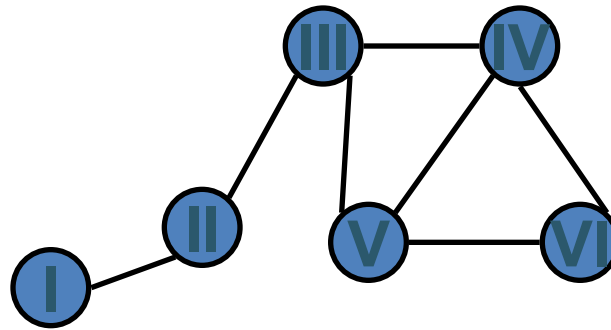
**X**

	I	II	III	IV	V	VI
I	0	1	0	0	0	0
II	1	0	1	0	0	0
III	0	1	0	1	1	0
IV	0	0	1	0	1	1
V	0	0	1	1	0	1
VI	0	0	0	1	1	0

**X<sup>2</sup>**

	I	II	III	IV	V	VI
I	1	0	1	0	0	0
II	0	2	0	1	1	0
III	1	0	3	1	1	2
IV	0	1	1	3	2	1
V	0	1	1	2	3	1
VI	0	0	2	1	1	2

# From Graphs to (Geodesic Distance)-Matrices (Reachability) – Geodesic Distance



Binary, undirected

	I	II	III	IV	V	VI
I	-	1	2	3	3	4
II	1	-	1	2	2	3
III	2	1	-	1	1	2
IV	3	2	1	-	1	1
V	3	2	1	1	-	1
VI	4	3	2	1	1	-

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