

# 6.035

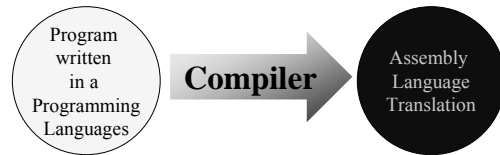
Fall 2005

## Lecture 15: Putting it all together

From parsing to code generation

## How to make the computer understand?

- Write a program using a programming language
- Microprocessors talk in assembly language



Saman Amarasinghe

2

6.035 ©MIT Fall 1998

## Example (input program)

```
int expr(int n)
{
    int d;
    d = 4 * n * n * (n + 1) * (n + 1);
    return d;
}
```

Saman Amarasinghe

3

6.035 ©MIT Fall 1998

## Example (Output assembly code)

### Unoptimized Code

```
expr:
.LFB2: pushq %rbp
.LCFI0: movq %rbp, %rbp
.LCFI1: movl %edi, -4(%rbp)
        movl -4(%rbp), %eax
        movl %eax, %edx
        imull -4(%rbp), %edx
        movl -4(%rbp), %eax
        incl %eax
        imull %eax, %edx
        movl -4(%rbp), %eax
        incl %eax
        imull %edx, %eax
        sall $2, %eax
        movl %eax, -8(%rbp)
        movl -8(%rbp), %eax
        leave
        ret
```

### Optimized Code

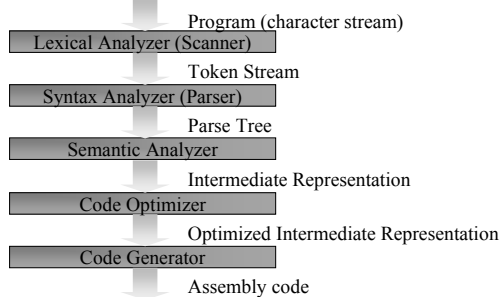
```
expr:
.LFB2: movl %edi, %eax
        imull %edi, %eax
        incl %edi
        imull %edi, %eax
        imull %edi, %eax
        sall $2, %eax
        ret
```

Saman Amarasinghe

4

6.035 ©MIT Fall 1998

## Anatomy of a Computer



Saman Amarasinghe

5

6.035 ©MIT Fall 1998

## Lexical Analysis

- Lexical analyzer create tokens out of a text stream
- Tokens are defined using regular expressions

Saman Amarasinghe

6

6.035 ©MIT Fall 1998

## Examples of Regular Expressions

Regular Expression	Strings matched
a	"a"
a · b	"ab"
a   b	"a" "b"
$\epsilon$	""
a*	"a" "aa" "aaa" ...
(a   $\epsilon$ ) · b	"ab" "b"
num = 0 1 2 3 4 5 6 7 8 9	"0" "1" "2" "3" ...
posint = num · num*	"8" "6035" ...
int = ( $\epsilon$   -) · posint	"-42" "1024" ...
real = int · ( $\epsilon$   ( · posint))	"-12.56" "12" "1.414" ...

Saman Amarasinghe

7

6.035 ©MIT Fall 1998

## Lexical Analysis

- Lexical analyzer create tokens out of a text stream
- Tokens are defined using regular expressions
- Regular expressions can be mapped to Nondeterministic Finite Automaton (NFA)
  - by simple construction
- NFA is transformed to a DFA
  - Transformation algorithm
  - Executing a DFA is straightforward

Saman Amarasinghe

8

6.035 ©MIT Fall 1998

## Syntax Analysis (parsing)

- Defining a language using context-free grammars

Saman Amarasinghe

9

6.035 ©MIT Fall 1998

## Example: A CFG for expressions

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle \langle \text{op} \rangle \langle \text{expr} \rangle$   
 $\langle \text{expr} \rangle \rightarrow ( \langle \text{expr} \rangle )$   
 $\langle \text{expr} \rangle \rightarrow - \langle \text{expr} \rangle$   
 $\langle \text{expr} \rangle \rightarrow \mathbf{num}$   
 $\langle \text{op} \rangle \rightarrow +$   
 $\langle \text{op} \rangle \rightarrow *$

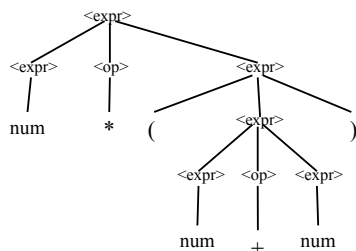
Saman Amarasinghe

10

6.035 ©MIT Fall 1998

## Parse Tree Example

num '\*' '(' num '+' num ')'



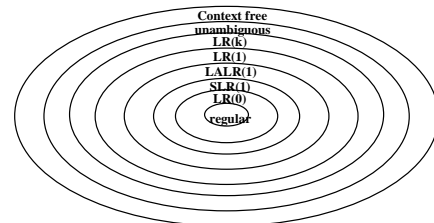
Saman Amarasinghe

11

6.035 ©MIT Fall 1998

## Syntax Analysis (parsing)

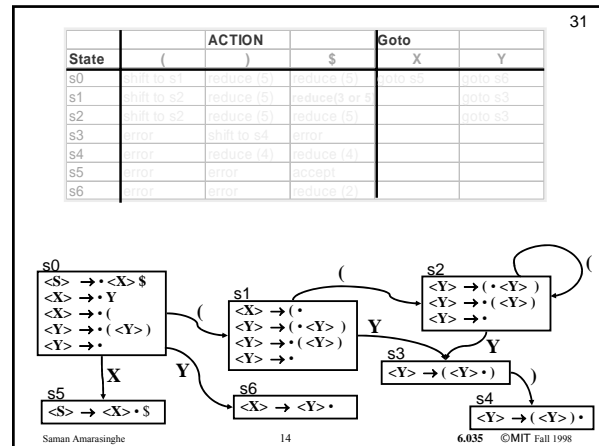
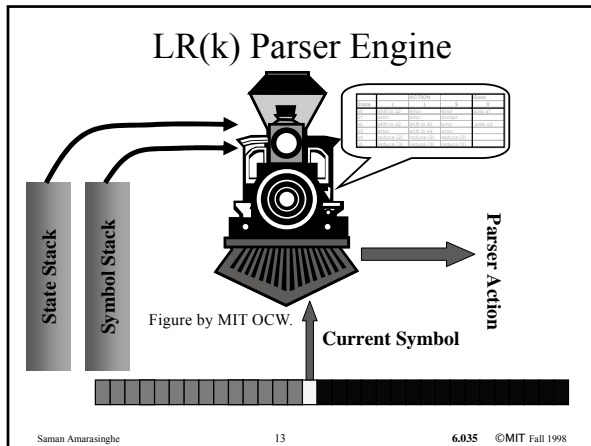
- Defining a language using context-free grammars
- Classification of Grammars



Saman Amarasinghe

12

6.035 ©MIT Fall 1998



- ### Semantic Analysis
- Building a symbol table
  - Static Checking
    - Flow-of-control checks
    - Uniqueness checks
    - Type checking
  - Dynamic Checking
    - Array bounds check
    - Null pointer dereference check
- Saman Amarasinghe
15
6.035 ©MIT Fall 1998

- ### Translation to Intermediate Format
- Goal: Remain Largely Machine Independent But Move Closer to Standard Machine Model
    - From high-level IR to a low-level IR
  - Eliminate Structured Flow of Control
  - Convert to a Flat Address Space
- Saman Amarasinghe
16
6.035 ©MIT Fall 1998

- ### Code Optimizations
- Generate code as good as hand-crafted by a good assembly programmer
  - Have stable, robust performance
  - Abstract the architecture away from the programmer
    - Exploit architectural strengths
    - Hide architectural weaknesses
- Saman Amarasinghe
17
6.035 ©MIT Fall 1998

- ### Code Optimizations
- Algebraic simplification
  - Common subexpression elimination
  - Copy propagation
  - Constant propagation
  - Dead-code elimination
  - Register allocation
  - Instruction Scheduling
- Saman Amarasinghe
18
6.035 ©MIT Fall 1998

## Compiler Project!

- You guys build a full-blown compiler from the ground up!!!
- From decaf to working code

## Compiler Derby

- Who has the fastest compiler in the east???
- Will give you the program 12 hours in advance
  - Test and make all the optimizations work
  - DO NOT ADD PROGRAM SPECIFIC HACKS!
- Wednesday, December 14<sup>th</sup> at 11:00AM  
location TBA
  - refreshments provided

## How will you use 6.035 knowledge?

- As an informed programmer
- As a designer of simple languages to aid other programming tasks
- As an engineer dealing with new computer architectures
- As a true compiler hacker

## 1. Informed Programmer

- Now you know what the compiler is doing
  - don't treat it as a black box
  - don't trust it to do the right thing!
- Implications
  - performance
  - debugging
  - correctness

## 1. Informed Programmer

- What did you learned in 6.035?
  - How optimizations work or why they did not work
  - How to read and understand optimized code

## 2. Language Extensions

- In many applications and systems, you may need to:
  - implement a simple language
    - handle input
    - define an interface
    - command and control
  - extend a language
    - add new functionality
    - modify semantics
    - help with optimizations

## 2. Language Extensions

- What you learned in 6.035
  - define tokens and languages using regular expressions and CFGs
  - use tools such as jlex, lex, javacup, yacc
  - build intermediate representations
  - perform simple transformations on the IR

## 3. Computer Architectures

29

- Many special purpose processors
  - in your cell phone, car engine, watch, etc. etc.
- Designing new architectures
- Adapting compiler back-ends for new architectures

## 3. Designing New Architectures

- Great advances in VLSI technology
  - very fast and small transistors
  - scaling up to billion transistors
  - but, slow and limited wires and I/O
- A computer architecture is a combination of hardware and compiler
  - need to know what a compiler can do and what hardware need to do
  - If compiler can do it don't waste hardware resources.

## 3. Designing New Architectures

- What did you learned in 6.035
  - Capabilities of a compiler: what is simple and what is hard to do
  - How to think like a compiler writer

## 3. Back-end support

- Every new architecture need a new backend
- Instruction scheduling
  - Even if the ISA is the same, different resource constrains
  - How to handle new features

## 3. Back-end support

- What do you learned in 6.035
  - Intermediate representations
  - Transforming/optimizing the IR
  - Process of generating assembly from a high-level IR
  - Assembly interface issues (eg: calling conventions)
  - Register allocation issues
  - Code scheduling issues

36

## 4. Compiler Hacking

- Theory
- Algorithms
- Implementation

Saman Amarasinghe 31 6.035 ©MIT Fall 1998

36

## 4. Compiler Hacking

- Theory:
  - Develop general, abstract concepts
  - Prove correctness, optimality etc.
- Examples
  - parse theory
  - lattices and data-flow
  - abstract interpretation
  - The language ML

Saman Amarasinghe 32 6.035 ©MIT Fall 1998

## 4. Compiler Hacking

- Algorithms:
  - Design a solution to a given problem (Mostly new optimizations)
  - Use many techniques such as graph theory, number theory, etc.
  - May have to limit the scope and find good heuristics
- Examples
  - partial redundancy elimination
  - register allocation by graph coloring
  - using multi-granular (MMX) operations

Saman Amarasinghe 33 6.035 ©MIT Fall 1998

## 4. Compiler Hacking

- Implementation:
  - Develop a new compiler
  - Issues of designing a very complex software
  - Putting theory and algorithms into practice
- Examples
  - A JIT for Java
  - A query optimization engine for SQL
  - A rasterizer for postscript

Saman Amarasinghe 34 6.035 ©MIT Fall 1998

## 4. Compiler Hacking

- What did you learned in 6.035?

*Everything!!!!*

Saman Amarasinghe 35 6.035 ©MIT Fall 1998

## Where to Look for Current Research?

- PLDI – Programming Languages Design and Implementation Conference
- Code Generation / Machine specific
  - Micro Conference
  - ASPLOS – Architecture Support for Programming Languages and Operating Systems
  - CGO – Code Generation and Optimization
- Language Theory
  - POPL – Principles of Programming Languages
  - OOPSLA – Object Oriented Programming Systems Languages and Applications
- Program Analysis
  - SAS – Static Analysis Symposium
  - PPOPP – Principles and Practice of Parallel Programming

Saman Amarasinghe 36 6.035 ©MIT Fall 1998